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**Patentanmeldung Nr.    Patent application No.    Demande de brevet n°**

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Der Präsident des Europäischen Patentamts;  
Im Auftrag

For the President of the European Patent Office

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-1-

**AUCTION METHOD AND APPARATUS FOR ELECTRONIC COMMERCE****Field of the Invention**

The present invention relates to computing entities configured for electronic  
5 commerce, and particularly although not exclusively, to a method and apparatus  
for making bids and offers at electronic auctions over the internet.

**Background to the Invention**

It is known on the internet that several commercial websites are present  
10 which serve the purpose of auctioning goods. Referring to Fig. 1 herein there is  
shown an example of a prior art system in which first and second commercial  
auction websites 100, 101 are each supported by a corresponding computer  
entity 102, 103 operating on a website server. A user of the auction websites  
over the internet 104 accesses the auction websites 100, 101 by use of his own  
15 computing entity 105. In Fig. 1, auction websites 100, 101 and the user are  
shown as logical entities (being websites supported by the computing entities)  
and the user operates a browser software 106 resident on the user's computer  
entity 105. Examples of such auction websites include those provided by eBay™  
and QXL. A human user having a conventional personal computer equipped with  
20 a conventional browser software may use the browser software to access the  
auction websites over the internet in known manner, and using the browser may  
generate a screen display downloaded from the website which provides the user  
with an interactive interface on his computer entity 105 through which the user  
can participate in an auction hosted by the website servers 102, 103. To  
25 participate in an auction, a user typically must register with the auction site server  
by entering the user's details such as company name, personal name, credit card  
number, and address, in to dialog boxes provided on the website interface. The  
user accesses information concerning goods sold through the auction website.

30990066

-2-

Referring to Fig. 2 herein, there is illustrated schematically aspects of a typical user interface display of a known auction website, which a human user can access on his or her graphical user interface by pointing a browser software at the website. Typically such a display comprises a listing of goods or services to be sold, a quantity of goods or services, a current highest bid price which is being bid for the goods or services, a number of bids which have been received by the website for those goods or services, a time at which the auction started, and a time at which the auctioning of the goods or services may close, as well as other details concerning payment method and shipping terms.

Referring to Fig. 3 herein there is illustrated an example of a prior art screen display accessed by a user for making a bid. A user enters an amount which the user bids for the goods using a dialog box provided. The bid is relayed over the internet to the website, which inputs the bid information into its controlling software to determine whether or not the bid is successful. Once the auction software accepts a successful bid, a confirmation message is sent to the computing entity of the successful bidder, and a transaction, including supply of goods and payment for goods, occurs by conventional means, for example postal delivery of goods and conventional credit card transaction clearance.

It is known in prior art systems to provide at a known auction website a software functionality known as a 'bidding elf', which is provided as part of an auction website, and which is capable of bidding in an auction on behalf of a buyer. The bidding elf provides a user with a screen display having dialog boxes into which the user can enter an upper maximum limit to which he or she is prepared to bid for various goods. The bidding elf then places bids on behalf of the user, varying the bid level incrementally up to the maximum predetermined limit set by the user. The bidding elf removes the need for the user to sit at his visual display continuously to monitor current bidding whilst the auction is occurring and continue to increase his bid incrementally up to his maximum bid.



30990066

-3-

With known electronic auctions, generally a user can only attend one auction at any one time, since the website browser can only be pointed at one auction at any one time and the user interface of one auction website can be  
5 displayed on a single monitor at any one time. If a user wishes to attend more than one auction simultaneously, in order to source goods or services at the best price from several auctions, then the user must arrange a plurality of website browsers supported by one or more computing entities to communicate with the plurality of auctions, and must visually monitor each separate auction site  
10 simultaneously. Such a system is possible, using conventional technology, but is inelegant, requiring several items of computing equipment by the user, and requiring manual monitoring and interaction of several auction sites at once. The prior art bidding elf can only make bids in one auction, being the auction website with which it is associated. Different prior art elves at different websites cannot  
15 communicate with each other to coordinate bidding in several auctions.

### Summary of the Invention

One object of the present invention is to provide a method and apparatus to enable a user to participate in several electronic auctions simultaneously in a  
20 controlled and coordinated manner to obtain a best price contract for goods and services from a plurality of auctions. In one implementation, a user may set an upper and/or lower price limit for bid or offer of goods or services at a plurality of auctions, and thereby participate in risk free buying or selling of goods and services. The user may be guaranteed that the maximum price bid will not be  
25 exceeded, or a minimum price offered will not be exceeded in specific embodiments of the invention.

Specific embodiments may enable a user to obtain items at a preferential rate in one auction whilst participating in one or more other auctions for the same  
30 goods/services and whilst avoiding buying of more goods than required, or selling

30990066

-4-

more goods than required through participating in other ones of the plurality of auctions.

According to first aspect of the present invention there is provided an  
5 electronic trading entity comprising:

a processor;

a memory means;

10

a communications means;

monitoring means for monitoring a plurality of trading data displayed by a  
plurality of auction entities;

15

quantity calculation means for calculating quantities of items for trading  
with said plurality of monitored auction entities,

price calculation means for calculating optimal prices of items to be traded  
20 with said plurality of auction entities;

wherein said price calculation means operates to process said trading  
data to obtain trading prices for a quantity of items calculated by said quantity  
calculation means, said prices calculated for a set of chosen auction entities  
25 selected from said set of all said plurality of auction entities;

said price calculation means operating to process said selected data to  
obtain trading prices for a predetermined amount of items across all said selected  
set of said plurality of auction entities.

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-5-

Preferably the electronic trading entity further comprises search means for searching electronically for individual ones of said auction entities.

Preferably said amount calculation means comprises:

5

means for storing bid data describing a plurality of bids in at least one of said plurality of auction entities;

means for storing user bid data describing a number of bids placed by said trading entity with at least one of said plurality of auction entities; and

10

means for determining an optimum number of bids and corresponding price amounts of said bids to place with at least one of said plurality of auction entities.

15

According to a second aspect of the present invention there is provided a method of operating an electronic trading entity for trading a plurality of tradable items, said method comprising the steps of:

20

monitoring a plurality of auction entities;

inputting a plurality of trading data from said plurality of auction entities;

25

processing said trading data of said plurality of auction entities to determine an optimum set of amount data describing quantities and prices of said tradable items for trading by said trading entity.

Said step of processing trading data of said plurality of auction entities may comprise:

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30990066

-6-

storing in memory a plurality of bid data displayed by said plurality of auction entities;

5 determining a number of active bids of said trading entity already sent by said trading entity to at least one of said plurality of auction entities;

constructing a plurality of currently successful sets of said bid data of said plurality of auction entities;

10 determining a cost to outbid each said set of currently successful bids;

selecting a said currently successful set of bids having an optimum cost to outbid; and

15 constructing a set of bids of said trading entity which are higher than corresponding bids of said optimum cost currently successful bid set.

Said step of constructing a set of bids of said trading entity which are higher than said optimum cost currently successful bid set may comprise:

20

for each bid of said optimum bid set, said bid made either by said trading entity or a third party, selecting a highest third party bid appearing in an identical auction entity to said bid, which also appears in said optimum bid set; and

25 generating a trading entity bid of an amount equal to said highest third party bid plus a minimum bid increment.

Said step of processing said trading data of said plurality of auction entities may comprise:

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30990066

-7-

storing in memory a plurality of offer data of said plurality of auction entities;

5 determining a number of active offers already placed by said plurality of auction entities;

constructing a plurality of sets of offers of said auction entities;

10 determining a value of each of said set of auction entity offers;

selecting a set of auction entity offers having an optimum value to undercut with new offers;

15 determining whether said value is within a predetermined value limit; and

generating a set of offer data matched to said optimum set of offer data.

Said step of processing data may comprise:

20 storing a plurality of offers of a plurality of auction entities,

selecting a plurality of sets of said plurality of offers, each set having a number of members equal to a number of items to be obtained by said trading entity;

25 for each said selected set, determining a value of said set;

selecting a said set having a maximum value;

30990066

-8-

generating a set of offers which correspond with said s.t having said maximum value.

Said step of generating a set of offers may comprise:

5

for each offer of said set of optimum offers, generating an offer having value equal to the corresponding respective offer value, minus a predetermined amount.

10

According to third aspect of the present invention there is provided a method of exchanging data between an electronic trading entity and a plurality of electronic auction entities, said method comprising the steps of:

15

said trading entity monitoring data displayed by said plurality of auction entities;

said first trading entity extracting from said monitored data, data describing prices of individual items to be traded;

20

said first trading entity determining from said price data an optimum set of trading data for sending to said plurality of auction entities; and

said first trading entity communicating said trading data to said plurality of auction entities.

25

### **Brief Description of the Drawings**

For a better understanding of the invention and to show how the same may be carried into effect, there will now be described by way of example only, specific embodiments, methods and processes according to the present invention with reference to the accompanying drawings in which:

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30990066

-9-

Fig. 1 illustrates schematically an example of a prior art system in which first and second commercial auction websites are each supported by corresponding computer entity;

5

Fig. 2 illustrates schematically an example of a typical prior art user interface display of a known auction website;

Fig. 3 illustrates schematically further features of the typical prior art user interface display of the known auction website;

10

Fig. 4 illustrates schematically operation of a trading entity according to first specific implementation of the present invention in an internet environment;

Fig. 5 illustrates schematically an architectural structure of the trading entity of Fig. 4;

15

Fig. 6 illustrates schematically interaction of the trading entity with a plurality of auction entities over the internet and interaction of the trading entity with a user;

20

Fig. 7 illustrates schematically a logical layout of the trading entity of Fig. 4;

Fig. 8 illustrates schematically process steps carried out by the trading entity for procuring a single good or item from a plurality of auction entities;

25

Fig. 9 illustrates schematically process steps carried out by a monitoring and trading component of the bidding algorithm for bidding for several goods or services at once from a plurality of auction entities;

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30990066

-10-

Fig. 10 illustrates schematically a data table recording data describing a plurality of bids and/or offers which the trading entity is actively pursuing at a plurality of auction entities;

5        Fig. 11 illustrates schematically a bid selection algorithm for selecting a set of bids to place across a plurality of auction entities;

Fig. 12 illustrates schematically an array of bids for each of a plurality of N auctions in which individual sets of bids D are selected;

10

Fig. 13 illustrates schematically an example of a specific set of bids selected exclusively from a single auction;

15        Fig. 14 illustrates schematically a set of bids comprising 4 bids, placed in 2 auctions at levels which outbid currently active bids of third parties in those auction entities;

20

Fig. 15 illustrates schematically an offer selection algorithm for selecting a set of offers to be made in a plurality of auction entities;

Fig. 16 illustrates schematically a method for determining an optimum set of bids to place with one or a plurality of auction entities;

25        Fig. 17 illustrates schematically a set of bids which may be displayed by a plurality of auction entities and stored by a trading entity according to a specific embodiment of the present invention; and

Fig. 18 illustrates schematically a set of offers which may be generated by the trading entity, to match with a set of bids as illustrated in Fig. 17 herein.

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30990066

-11-

**Detailed Description of the Best Mode for Carrying Out the Invention**

There will now be described by way of example the best mode contemplated by the inventors for carrying out the invention. In the following description numerous specific details are set forth in order to provide a thorough understanding of the present invention. It will be apparent however, to one skilled in the art, that the present invention may be practiced without limitation to these specific details. In other instances, well known methods and structures have not been described in detail so as not to unnecessarily obscure the present invention.

10

A first specific implementation according to the present invention is concerned primarily with electronic commerce according to the 'English auction' system in which a seller takes competing bids from a plurality of bidders, with a current highest bid price from one of the bidders being the price which other bidders must exceed in order for a successful transaction to occur. Where the highest bid price is not exceeded for a predetermined period, the seller may choose to either withdraw the goods from auction if his reserve price is not met, or the seller may enter into a transaction with the bidder making the highest bid in the auction.

20

Throughout this description, transactions concerning goods or services will be referred to. Where a transaction to buy or sell a good is mentioned, it will be understood that the subject of the transaction may be either a good or a service or both a good and service, the terms good and service being used interchangeably, and similarly a reference to a transaction involving a service will be construed accordingly to include a transaction involving a good.

25

In this specification the term "lot" is used to refer to a minimum quantity which is subject to a single transaction. A number of individual goods or services comprising a lot may be defined by a user of a computer trading entity according

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30990066

-12-

to specific implementations of the present invention described herein or may be specified by an entity with whom the computer trading entity trades. A user may define a lot size for any particular type of good or service as being one, or a preset plurality of goods or services.

5

Referring to Fig. 4 herein, there is illustrated schematically an interaction of a trading device 400 according to a first specific implementation of the present invention with a plurality of auction entities 402 – 404. The trading device comprises a computer program operating to control a first computing entity 401  
10 comprising a data processor and memory means. Each auction entity comprises a computer program operating a corresponding computing entity 405 – 407, the computing entities each comprising a data processor and associated memory means, the trading device 400 communicating with the one or plurality of auction entities over the internet 408.

15

Referring to Fig. 5 herein, there is illustrated schematically an architectural structure of trading device 400. The trading device comprises a data processing means 500; an area of memory 501 for storage of data and instructions used by processor 500; a user interface 502, comprising a video display screen, keyboard  
20 and pointing device, for example a mouse device; an operating system 503, for example the known Unix™ operating system, Microsoft Windows™, operating system, or Linux™ operating system; a modem 504 for communicating with other internet devices; a plurality of applications 505, including a web browser 506, and a set of algorithms 507 for placing bids and offers at auction sites, the algorithms  
25 507 co-operating with web browser 506. The web browser 506 may comprise a conventional web browser, for example the known Microsoft Internet Explorer™ browser.

Referring to Fig. 6 herein, there is shown schematically communication of  
30 trading device 400 with the plurality of auction entities, and with a user of the

30990066

-13-

device. Trading device 400 receives hypertext mark-up language (HTML) or XML signals 600 from the auction entities, and sends data 601, also in HTML or XML to the auction entities over the internet. Trading device 400 communicates with the user interface using HTML signals, and can receive data input and commands from a user via the user interface, for example by the user typing in commands, or using the pointing device.

The trading device comprises a local set of software commands on the computer entity 401, which operate to parse information from the auction entities, and receives HTML signals, in the same way as the web browser. The trading device extracts information from the HTML signals concerning a current status of an auction entity, including information on a highest current bid and the bidding increments set by the auction entity.

The trading device 400 contains a plurality of algorithms for interacting with the user and one or more auction entities as will now be described.

Referring to Fig. 7 herein, there is illustrated a logical layout of trading device 400 including arrows indicating flows of data between logical components of the trading device 400. Trading device 400 comprises a web browser 506; monitoring and trading algorithms 700; one or more price calculation algorithms 701, an auction selection algorithm 702, and a URL store 703. Web browser 506 may monitor one or a plurality of auction entities under control of monitoring and trading algorithm 700. Monitoring and trading algorithm 700 interacts with price calculation algorithm 701 to obtain information concerning calculated optimum prices and lowest/highest prices. Price calculation algorithm 701 communicates with auction selection algorithm 702 for selecting auction entities having the best prices for selected goods or services. URL store 703 stores URL addresses of selected auction entities which are referred to by the auction selection algorithm 702 and monitoring and trading algorithm 700. User interface 502 communicates

30990066

-14-

with web browser 506, and monitoring and trading algorithm 700, so that a human user of the trading device can monitor and participate in trades carried out by the monitoring and trading algorithm 700 and actively search and select auction entities using web browser 506. Each of monitoring and trading algorithm  
5 700 and web browser 506 are capable of communicating with auction entities via the modem facility 504 of the host computer entity, via the internet.

Whilst communications between the trading device and the auction entities are described herein as being carried out in HTML or XML, protocols operated by  
10 the trading device 400 for communication in the general sense is not limited to these languages, and in the general case communication may occur by other standard interfaces, for example FIPA (foundation of intelligent physical agents) defined languages.

15 In some embodiments, algorithms for communicating with agent entities may reside within a web browser, in which case the algorithms may generate HTML driven screens which the user may review using his conventional web browser. Information which the trading device displays to the user includes information concerning the status of a set of auction sites selected by the user.  
20 The trading device may be directed to web sites by a user typing in addresses of known websites, using the URL of the websites, or alternatively a communication algorithm within the device may be configured to use the browser to search for auction sites by entering in a keyword search via the browser, into a conventional search engine, e.g. Yahoo™ or Alta Vista™.

25

A user of the trading device needs to have confidence in the integrity and reliability of the auction entities which the trading device is communicating with. Consequently, in the best mode implementation, rather than leave selection of auction entities to the trading device, a user may select a plurality of auction  
30 entities with which the trading device communicates, for example by inputting the

30990066

-15-

URL data of selected auction entities to the trading device, which the trading device stores in URL memory store 703. Other user inputs to the trading device include:

- 5           •     A list of goods or services which the user wishes to buy, sell or monitor for. Goods need not necessarily be actively bought or sold, but can be monitored and a user alerted to any particular goods of interest when these arise on an auction site.

- 10           •     Information concerning auctions which the user wishes to participate in. This information may include a general instruction to participate in any auction which the user has pre-selected and has entered its URL into the trading device.

- 15           •     A quantity of goods or services to buy or sell. This data may comprise a maximum quantity of goods or services to buy or sell, and a minimum quantity of goods or services to buy or sell. For example, when buying or selling electronic components, lots of 1000 units are common, and lots of 100 units or lower are generally not traded.

- 20           •     In the case of a buy instruction from the user to the trading device, information describing an upper bid price limit (and currency type) which the user is prepared to pay without the trading device referring back to the user.

- 25           •     In the case of buying goods or services, an absolute upper bid price limit and currency above which the user is unwilling to purchase the goods or services.

- 30           •     In the case of selling goods or services, a lower offer limit below which the user requires referral of price information before selling the goods or services.

30990066

-16-

- In the case of selling goods or services, an absolute lower offer limit below which the user is unwilling to sell the goods or services.

If the user is buying, and the goods are within the user's upper bid limit, the trading device may buy goods and services automatically without further referral to the user. Similarly, if the user is selling goods and services and the prices obtained are above the user's lowest offer limit, the trading device may sell goods and services automatically without further referral to the user.

Referring to Fig. 8 herein, there is illustrated steps of a simple algorithm for purchasing one lot of goods or service only from one of a plurality of N auction entities.

In step 800, the user inputs a maximum price limit  $P_{MAX}$  being a price which should never be exceeded in any of the N auctions for a quantity of goods or services.

In step 801, the user inputs a maximum price data  $P_{MAX}$  into his user interface, which is received by trading device 400. In step 801, the user enters a referral price  $P_{REF}$ , which is a price above which any bids to be made by the device must be referred back to the user before proceeding. In step 802, the trading device monitors a plurality N of pre-selected auction entities, which have previously been selected by the user, and which are selling the goods or services of interest to the user. In step 803, the trading device monitors the current highest bids  $b$  in each of the plurality of N auctions monitored. This results in a plurality of bids  $b_1, \dots, b_N$ , being the current highest bid operating in each of the N auctions. In step 804, if the trading device has a currently entered bid in any of the N auctions, which is the highest bid in that auction, that is to say if the trading device is actually leading the bidding in any one of the plurality of N auctions, then the algorithm idles, and continues to monitor the N auctions and current bid

30990066

-17-

prices in steps 802, 803 as described previously. However, if the trading device is not currently leading the bidding in any one of the  $N$  auctions, in step 805 the algorithm selects the auction having the lowest current active bid (the active bid being the highest bid in that auction) in step 805. In step 806, it is checked  
5 whether there is more than one auction having an equivalent lowest active bid. If there is only one such auction, then in step 809, the algorithm calculates a bid at the level  $b_i + \delta$  in the selected auction  $i$ . However, if there is more than one auction having a lowest active bid price, that is to say that there is a same lowest active bid price for the goods or services in more than one auction, then the  
10 algorithm selects from these the auction having the earliest finish time, and the lowest active bid price in step 807. In step 808, the algorithm then proceeds to calculate a bid  $b_i + \delta$  in the selected auction  $i$ , where  $\delta$  is the smallest bid increment acceptable in that auction  $i$ . In step 810 the algorithm compares the bid  $b_i + \delta$  with the price limit above which all bids must be referred to the user  
15  $P_{REF}$ . If the bid  $b_i + \delta$  is below the referral limit  $P_{REF}$ , in step 811, the algorithm proceeds to place the bid  $b_i + \delta$  in the selected auction and returns to monitoring state 802. However, if the bid  $b_i + \delta$  is above the limit  $P_{REF}$  above which all bids must be referred to the user, in step 812, the algorithm compares the bid  $b_i + \delta$  with the absolute maximum price limit  $P_{MAX}$  set by the user to see if it is below  
20  $P_{MAX}$ . If the calculated bid exceeds  $P_{MAX}$  then the algorithm stops and no further bidding takes place. The algorithm stops in step 813. However, if the calculated bid  $b_i + \delta$  is within the maximum upper price limit  $P_{MAX}$ , in step 814, the algorithm refers the bid to the user for authorization. If the user authorizes the bid in step 815, then the bid  $b_i + \delta$  is placed in the selected auction in step 816 and returns  
25 to monitoring step 802. However, if the user does not authorize the bid, or rejects the bid  $b_i + \delta$  then in step 817, the algorithm stops and no further bidding takes place.

The above algorithm ensures that the user always has a leading bid in one  
30 of the plurality of  $N$  auctions, placing the lowest possible bid to take a lead in one

30990066

-18-

of the plurality N auctions and without placing a bid in any other auction than the one in which the lowest possible bid is placed until that bid is overtaken. When the user's bid is overtaken in the auction in which the user is bidding, the algorithm selects the next auction having the lowest active bid price and places a  
5 bid in that auction. This auction may be the same auction as the one in which the user has just been outbid in, or may be a different auction of the plurality N auctions. The algorithm is always searching for the lowest current active bid which allows the user to lead in one of the plurality N auctions for the required single lot of goods or services, and always below the user's referral price limit  
10  $P_{REF}$  which the user is prepared to pay without the trading device referring back to the user.

The referral mechanism provided where a bid is above the user's price limit above which all bids need to be referred back to the user  $P_{REF}$ , but below the  
15 user's absolute maximum price limit  $P_{MAX}$  above which the user will never want to bid for the lot of goods or services, is provided primarily to give the user confidence in the trading device and gain acceptance of the bidding system from the user. However, an optimally behaving user should not use the referral mechanism, but should set his maximum price limit above which he would never  
20 wish to pay to be the same as the price limit above which referrals should be made, that is to say  $P_{REF} = P_{MAX}$ . The trading device will never pay more than the upper limit set by the user, and will not reveal that upper limit to the auction, or to competing bidders.

25 Referring to Fig. 9 herein, there is illustrated schematically an example of procedure steps carried out by monitoring and trading algorithm 100 for bidding for several lots of goods or services at once at a plurality of auction entities. A number M goods of a specified type are to be purchased from a plurality of N auctions at an optimal overall price, that is to say, the best prices available  
30 throughout the N auctions. Using this algorithm, some of the plurality M goods



30990066

-19-

may be bought at different prices to others, but overall, the best prices for the M goods are obtained, subject to availability, at the plurality N of auctions. In each of the auctions, different bidders may bid a different price per good for different lots of goods where the price depends on the quantity traded. For example, a first bidder may bid say \$500 each for 5 items (5 lots) whereas in another auction, another bidder may bid \$450 per item for 10 items (10 lots) of the same goods. The algorithm of Fig. 9 deals with the feature of the bid price per lot being dependant upon the size of transaction by treating each of the M lots to be traded separately, and attaching a separate bid price to each one of the plurality M lots to be bid for. Therefore, in the algorithm, instead of bidding \$2,500 for 5 lots, the algorithm breaks the transaction down into 5 individual bids of \$500 for each one of 5 lots, resulting in 5 bids instead of 1 bid in this example.

In step 900, the user selects each of the plurality of N auctions in which the user wishes to bid for the goods or services. This may be achieved by selecting the auctions using a conventional web browser and inputting the URL of each selected auction into the trading device. In step 901, the user inputs the single type of good or service, and the total number of lots M of the good or service to be bid for in the plurality N auctions via the user interface 502. In step 902, the user inputs an absolute maximum bid price  $P_{MAX}$  per lot, which the user is prepared to pay. A lot of goods comprises a predetermined quantity of goods or services which the user is buying, and may be a single good, or a plurality of goods, for example 1000 individual items. In step 903, the user inputs a referral price,  $P_{REF}$  being a cost amount per lot above which any bids must be referred to the user before they can be entered at any of the auctions. As described above, a user who behaves optimally will set  $P_{REF} = P_{MAX}$  so that in practice no referral of price occurs back to the user. In step 904, each one of the plurality of N auctions are continuously monitored to check the current active bid prices, where the active bid price is the highest bid price currently operating by any bidder in that individual auction. In general, the current active bid price in each of the plurality

30990066

-20-

of auctions will differ, but in a theoretically 100% efficient marketplace, the currently active bid price in all auctions may be the same. In step 905, for each auction  $i$  the algorithm considers the top  $M_i$  bids

5         $b_1^i, \dots, b_{M_i}^i$   
          where  $b_1^i \geq b_{M_i}^i$

For each auction, the top  $M_i$  bids are stored in a table, with the highest bid price (the highest currently active bid price which would be successful in any  
 10 auction)  $b_1^i$  at the top of the data table and the lowest active bid price  $b_{M_i}^i$  at the bottom of the table,  $b_{M_i}^i$  being the lowest bid price in the auction which would succeed in buying the goods or services. Further bids in the auction, which are currently unsuccessful (inactive) are labeled  $b_{m_i+1}^i, b_{m_i+2}^i \dots$  and so on.

15        The monitoring and trading algorithm maintains a database in memory in which for each of the  $N$  auctions which are being monitored, there is stored data describing a price amount for a plurality of individual bids which the algorithm has placed and which are currently active. An active bid is defined herein as a bid which has been placed with an auction entity, and which is of an amount  
 20 sufficient to secure a trade for that lot of goods or services in the absence of any subsequent higher bids for those same goods or services, and disregarding whether the amount exceeds a reserve price limit set by the auction entity for that same lot of goods or services.

25        The algorithm may also store connected with the active bid information for each lot, further information describing the proposed trade, including for example a description of the goods or services for which the bid is being made, a quantity of individual items for which the bid is being made, that is to say the number of individual items in a lot which is being bid for, a close time by which the bid will be  
 30 accepted or rejected by the auction entity, and a quality information describing a

30990066

-21-

quality of goods or services to be provided. Examples of the types of information which may be stored by the algorithm are illustrated in Fig. 10 herein. The data stored may comprise a bid identifier data 1000 being a unique identification code generated by the bid algorithm for a particular bid which has been placed; an  
5 auction identification data 1001, comprising a URL of the auction in which the bid identified by corresponding respective bid identifier 1000 has been placed; a lot identification data 1002, indicating a specific lot of goods or services for which the bid has been placed; a description information 1003 comprising a brief description of the goods or services for which the bid has been placed, or  
10 alternatively comprising a pointer to a separate table in which a more detailed description may be stored; a lot size, comprising the number of individual items (goods or services) in the lot which has been bid for; a monetary amount bid 1005, including a currency indicator relating to the currency in which the bid has been made, a time and date data 1006 at which the bid was made; a closing time  
15 of the auction for the particular lot which has been bid for; and a status indicator 1008 describing whether the current bid is in an active or inactive state.

In step 907, the monitoring and trading algorithm sums the number of bids, being one bid per specified lot, for the specified goods/services of the same type  
20 currently active, in all the selected auctions. For example, referring to Fig. 10, there are 5 bids active for lots of ten thousand 10k resistors, all of which are currently active, and spanning two individual auctions. In this example, the lot size is ten thousand individual resistors, and the price per lot which has been bid varies from \$25 to \$28. If  $L_i$  is the number of bids currently active which have  
25 been placed by the bidding device (in the example of Fig. 10 for the 10k resistor lots),  $L_i$  equals 5, then if for all auctions  $i = 1$  to  $N$ , then if the total number of bids  $L_i$  is greater than or equal to  $M$  as determined in step 908, then no further bids are placed and the algorithm continues to monitor each of the plurality of  $N$  auctions in step 904 and continues with subsequent steps 905 – 907. However, in step

30990066

-22-

908 the number of current active bids for all auctions is less than the total number of lots of goods or services to be obtained, that is to say:

Equation 1

5

$$\sum_{i=1}^N L_i \leq M$$

10 Then, the algorithm places additional bids according to a bid selection algorithm hereinafter described with reference to Fig. 11, so that an additional Q active bids are held by the trading entity, where Q is determined in equation 2 below.

15 Referring to Fig. 11 herein, there is illustrated an auction selection algorithm for selecting a subset of the plurality N of selected auctions in which to place new bids. The algorithm in Fig. 11 operates in real time, continuously inputting new bid values from each of the plurality of N auctions as they occur, and placing new bids on an ongoing basis as bids which have previously been placed by the algorithm are overridden by competing bids from third parties.

20

25 An object of the algorithm is to determine the lowest cost bids to place in the plurality of auctions N, in order to obtain the required number of lots of goods or services 1100. In step 1101, the algorithm receives input of data describing the current active bids of all of the plurality of N selected auctions. These bids comprise all the bids which, in the absence of any new bids, will be successful in obtaining goods or services (subject to these bids not being below a reserve price set by the auction). These bids are shown schematically in Fig. 12 herein as an array of bids for each auction i of the plurality N auctions. For the purpose of clarity in this example diagram, we assume every auction is selling exactly 4 lots.

30

Equation 2

30990066

-23-

$$Q = M - \sum_{i=1}^N L_i \leq M$$

- 5           The trading device then determines the optimum least cost bids to place which will result in an additional Q active bids.

In each auction A, there are a plurality of bids which need to be beaten in order to be successful in acquiring a given number of goods j from that auction.

10   A plurality of potential sets of bids are constructed to outbid existing bids in the plurality N of auctions as illustrated in step 1102 herein. The beatable – j set of existing bids in auction A is defined as the set

Equation 3

15            $(b_{Mi}^j, \dots, b_{Mi}^j)$  for  $j \geq 1$

This is the set of number j lowest bids which would succeed if the auction were terminated immediately. For example, in Fig. 12, 1201 shows the beatable – 1 set for auction 1, and 1202 shows the beatable – 2 set for auction 3.

20

The beatable zero (beatable – 0) set of bids for any auction is defined to be an empty set of bids.

A plurality of sets of bids can be formed by taking the members of exactly one beatable – j set ( $j \geq 0$ ) from every auction  $A_i$ . The set of all sets which can be formed by taking the members of exactly one beatable – j set ( $j \geq 0$ ) from every auction  $A_i$  is defined as the set  $\Phi$ . In Fig. 12 there are illustrated 3 members of  $\Phi$ , each with 3 elements;

25

- 30           • The set made from the beatable – 1 set of auction 1 (1202), and the beatable – 2 set of auction 3, and the beatable - 0 sets of all other auctions.

30990066

-24-

• The set (1203) made of the beatable – 3 set of auction N, and the beatable – 0 sets of all other auctions.

- 5       • The set (1204) made up of the beatable – 1 sets of auctions 1, 2 and 3 and the beatable – 0 set of all other auctions.

All sets D in the set  $\Phi$  which do not contain exactly Q bids made by parties other than the trading device are deleted.

10

(1103) - For each set  $D \in \Phi$  a cost of the set D is defined

Equation 4

$$15 \quad \text{Cost}(D) = \sum j_i (b_{mi}^{i-j_i+1} + \delta_i) - \sum_{i=1}^N b_{b \in D_A}$$

Where  $D_A$  is the set of bids in set D made by the trading device and  $j_i$  is the number of the highest bid which has been chosen by the algorithm in a particular set of bids of a selected auction  $A_i$ .

20

Equation 4 calculates the cost of outbidding each beatable – j set in D by placing  $j_i$  new bids of value  $b_{mi}^{i-j_i+1} + \delta_i$ , and subtracts the cost of existing bids in these sets made previously by the bidding device. This is illustrated schematically in Fig. 13, for a simple case where  $Q = 4$ , and the set D comprises the highest 4 bids in a single auction. In order to be successful in auction M, 4 bids 1302 must be placed, which exceed the highest bid  $b_1^M$  in auction M. The cost of placing these bids =  $4 (b_1^M + \delta)$ , where  $\delta$  is the minimum bid increment allowable in the auction. Similarly, in another example illustrated schematically in Fig. 14 herein, there is shown another set  $D_{pr}$  comprising a number  $Q = 4$  bids, placed in order to outbid the currently active bids of third parties in auctions P and

25

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30990066

-25-

R. To be successful in auctions P and R, 4 bids need to be placed 1501, 1502, being 2 bids of an amount  $b^P_3 + \delta$  in first auction P, and 2 bids each at an amount  $b^R_3 + \delta$  in auction R, giving a total cost  $2(b^P_3 + b^R_3 + \delta_P + \delta_R)$  for the bid set  $D_{PR}$ , where  $b^P_3$  is the second lowest currently active bid price in first auction P,  $\delta_P$  is the minimum bid increment in first auction P,  $b^R_3$  is the second lowest currently active bid price in second auction R, and  $\delta_R$  is the minimum bid increment in second auction R.

$\delta_i$  is the bid increment, which is specified by a particular auction. Bid increment data is obtained by a parser in the monitoring algorithm which reads the bid data directly from the auction entity, the data being transferred in a language such as HTML or XML.

(1104) - A set of bids having a minimum cost  $D_{MIN}$  is obtained by choosing  $D_{MIN} \in C$  such that

Equation 5

$$\text{Cost}(D_{MIN}) \leq \text{Cost}(D)$$

$$\forall D \in \Phi$$

If more than one set satisfies equation 5, then the latest auction finish times in each set are considered. The set with the earliest such finish time is chosen.

In step 1105 the algorithm then proceeds to deploy those bids in the relevant auctions at the appropriate calculated levels, as long as those bids are below the referral price  $P_{REF}$  set by the user. Any new bids which are calculated by the algorithm which are above the referral price  $P_{REF}$  will be referred back to the user. Under no circumstances will new bids be placed where the amount is

30990066

-26-

greater than the maximum price limit  $P_{MAX}$ . If all bids are placed, this will result in  $j_i$  bids of amount  $b_{mi-j}^i + \delta_i$  being placed in each auction  $A_i$ .

Referring to Fig. 15 herein, there is illustrated schematically an offer selection algorithm for selecting a set of offers to sell a plurality of lots of goods or services in a plurality of various auctions purchasing goods or services from the trading device. The trading device attempts to sell goods by making offers to a competitive tender. Hence the 'auctioneer' is putting out goods to competitive tender. The auctioneer will never receive bids, only offers.

The offer selection algorithm illustrated in Fig. 15 operates where a plurality N of auction entities request goods or services sold by the trading device 400 which offers those goods and/or services for sale. Each individual auction  $i$  buys a quantity  $M_i$  lots of goods or services. That is to say each auction entity  $i$  posts on its website details of the quantity of lots of goods or services which it wishes to acquire, and posts a call for offers for each of those lots of goods or services.

The trading device 400 places a plurality of offers with a selection of a plurality of N auction entities. The number of individual offers, the particular auction with which the offer is placed, and the amount of the offer are determined by the offer selection algorithm illustrated schematically in Fig. 15 herein.

The trading device monitors the plurality N auctions, each auction entity wishing to acquire a quantity  $M_i$  goods, the number  $M_i$  varying from auction to auction in the general case.

In step 1500, the user inputs details defining the quantity of goods or services M required to be sold or supplied into the trading device 400 using the user interface. These details are stored in local memory in the trading device for



30990066

-27-

further reference. In step 1501, the user inputs an absolute price limit  $P_{\text{MIN}}$  for each lot of goods or services which are to be sold, the absolute price limit  $P_{\text{MIN}}$  being the price below which the user is unwilling to sell or supply the goods or services. A separate price limit  $P_{\text{MIN}}$  is input for each lot of goods or services to be sold or supplied. The data describing the minimum offer price  $P_{\text{MIN}}$  is stored in memory by the trading device. In step 1502, the user inputs a referral price data  $P_{\text{REF}}$ , for each lot of goods or services which are to be sold or supplied. The referral price is the price below which the trading device cannot operate automatically, but must refer all offers which it makes to the plurality N of auctions back to the user via the user interface, before proceeding to place those offers. A user behaving optimally will set the minimum price  $P_{\text{MIN}}$  to be the same as the referral price  $P_{\text{REF}}$ , so that in practice the trading device once instructed to sell items does not refer back to the user. In step 1503, the trading device monitors the plurality of N auctions, by parsing data from the websites operated by those auction entities, the data describing the quantity, and offer price for goods or services of the type specified in step 1500. Auction entities may be selected on the basis of reliability, settlement terms, delivery terms, and trading history. In the best mode implementation, a user will personally select the plurality of N auction entities from which the price calculation algorithm can choose, rather than leaving the selection process to a further set of auction selection algorithms 702.

In step 1504, the algorithm monitors a plurality of offers from each of the plurality N auctions, receiving the lowest current offers in each auction i of the plurality. Where a large number of offers are posted by an auction entity, the algorithm may be configured to only receive from each auction at most a number of lowest offers  $S_1^i, \dots, S_{m_i}^i$  being the same as the quantity Q of the lots of goods/services to be disposed of. This prevents the trading device becoming flooded with large quantities of offer data from the plurality of auction entities. Auction entities may be pre-selected by the user in a similar fashion as described herein above with respect to an operation by the trading device.

30990066

-28-

In step 1504 for each auction  $i$  the lowest  $M_i$  offers  $S_1^i, \dots, S_{m_i}^i$  are stored. In step 1505, for each auction, the current number of active placed offers which the trading device has placed in the auction are stored. In step 1506, for all  
 5 selected auctions, the number of active offers of goods/services specified are summed. In step 1507, if the number of current active offers is greater than  $M$ , then the algorithm returns to step 1503 and the monitoring process is continued. However, in step 1507, if the number of current active offers is less than  $M$ , then in step 1508 the algorithm proceeds to place a set of new offers according to the  
 10 offer selection algorithm.

Selection of individual auctions with which to trade are chosen from the selected plurality  $N$  of auctions, and the prices at which those trades occur is carried out by the offer selection algorithm illustrated with reference to Fig. 16  
 15 herein. The basis for operation of the algorithm of Fig. 16 is as follows.

In step 1601, the plurality  $N$  of pre-selected auctions are monitored by parsing data from the websites of those auction entities. Offer data  $S_1^i, \dots, S_N^i$ , for a plurality of individual offers  $S$  for each auction  $A_i$  up to the maximum number of  
 20 auctions  $N$  is received by the trading device, suitably in HTML, or XML languages.

For an auction  $A_i$ , a beatable –  $j$  set of offers is defined as being:

25  $(S_{m_i-j+1}^i, \dots, S_{m_i}^i)$  for  $j \geq 1$ .

The beatable – 0 set for any auction is defined to be the empty set.

30990066

-29-

In step 1602, there are generated all sets of offers which can be formed, by taking the members of exactly one beatable – j set ( $j \geq 0$ ) from every auction  $A_i$ . Let  $\Phi$  be the set of all such sets.

5 All sets in  $\Phi$  which do not contain a number exactly Q of offers made by parties other than the trading device are deleted or ignored in step 1603. The remaining offer sets each having Q members are selected in step 1604.

For each set,  $D \in \Phi$  a profit D is defined as

10

Equation 6

$$\text{Profit}(D) = \sum_{i=1}^N j_i (S_{m_i-j_i+1} + \delta) - \sum_{S \in D_A} S$$

15

Where  $D_A$  is the set of offers in set D made by the trading device.

In step 1605, a maximum offer set  $D_{MAX}$  is selected by choosing  $D_{MAX} \in \Phi$   
20 such that

$$\text{Profit}(D_{MAX}) \geq \text{profit}(D)$$

for all  $D \in \Phi$ .

25 In step 1606 offers are deployed in the plurality of auctions by placing  $j_i$  offers each of amount

$$S_{m_i-j_i+1} - \delta$$

30 in each auction  $A_i$ , where  $\delta$  is an offer increment selected by the user or specified by the auction site.

30990066

-30-

Once the offers are placed with the selected auctions, then provided no competing offers are made by third parties at a lower price than those made by the trading device, then the offers should be accepted by the selected auctions.

5

Different numbers of lots of goods or services may be accepted by different auction entities, possibly at different times. Thus, the algorithm of Fig. 16 may dynamically and continuously loop back to step 1601 further monitoring the plurality of N auctions, reviewing any offers which have not yet been accepted. If an offer is accepted, then for that particular offer which is accepted, settlement occurs in step 1608, by the algorithm confirming by a message to the auction entity concerned that the offer has been accepted. This message forms a legally binding contract between the trading device and the auction entity. Payment and delivery may occur by conventional means, for example physically shipping the goods, and making a credit card transaction, or electronic funds transfer by conventional means. In the best mode implementation described herein, settlement of trades may be delegated to a human user for actual implementation.

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30990066

-31-

**Claims:**

1. An electronic trading entity comprising:

a processor;

5

a memory means;

a communications means;

10 monitoring means for monitoring a plurality of trading data displayed by a plurality of auction entities;

quantity calculation means for calculating quantities of items for trading with said plurality of monitored auction entities,

15

price calculation means for calculating optimal prices of items to be traded with said plurality of auction entities;

20 wherein said price calculation means operates to process said trading data to obtain trading prices for a quantity of items calculated by said quantity calculation means, said prices calculated for a set of chosen auction entities selected from said set of all said plurality of auction entities;

25 said price calculation means operating to process said selected data to obtain trading price for a predetermined amount of items across all said selected set of said plurality of auction entities.

30 2. The electronic trading entity as claimed in claim 1, further comprising search means for searching electronically for individual ones of said auction entities.

30990066

-32-

3. The trading entity as claimed in claim 1, wherein said amount calculation means comprises:

5 means for storing bid data describing a plurality of bids in at least one of said plurality of auction entities;

means for storing user bid data describing a number of bids placed by said trading entity with at least one of said plurality of auction entities; and

10

means for determining an optimum number of bids and corresponding price amounts of said bids to place with at least one of said plurality of auction entities.

15 4. The trading entity as claimed in claim 1, wherein said quantity calculation means comprises:

means for storing an offer data describing a plurality of offers made by said plurality of auction entities;

20

means for storing user offer data describing a number of offers placed by said trading entity with at least one of said plurality of auction entities; and

25 means for determining an optimum number of offers and corresponding price amounts of said offers to place with at least one of said plurality of auction entities.

5. A method of operating an electronic trading entity for trading a plurality of tradable items, said method comprising the steps of:

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30990066

-33-

monitoring a plurality of auction entities;

inputting a plurality of trading data from said plurality of auction entities;

5           processing said trading data of said plurality of auction entities to determine an optimum set of amount data describing quantities and prices of said tradable items for trading by said trading entity.

6.       The method as claimed in claim 5, wherein said step of processing  
10       trading data of said plurality of auction entities comprises:

          storing in memory a plurality of bid data displayed by said plurality of auction entities;

15           determining a number of active bids of said trading entity already sent by said trading entity to at least one of said plurality of auction entities;

          constructing a plurality of currently successful sets of said bid data of said plurality of auction entities;

20

          determining a cost to outbid each said set of currently successful bids;

          selecting a said currently successful set of bids having an optimum cost to outbid; and

25

          constructing a set of bids of said trading entity which minimally outbid said optimum cost to outbid currently successful bid set.

30990066

-34-

7. The method as claimed in claim 5, wherein said step of constructing a set of bids of said trading device which are higher than said optimum cost currently successful bid set comprises:

5 for each bid of said optimum bid set, said bid made either by said trading entity or a third party, selecting a highest third party bid appearing in an identical auction entity to said bid, which also appears in said optimum bid set; and

generating a trading device bid of an amount equal to said highest third  
10 party bid plus a minimum bid increment.

8. The method as claimed in claim 5, wherein said step of processing trading data of said plurality of auction entities comprises:

15 storing in memory a plurality of offer data of said plurality of auction entities;

determining a number of active offers of said trading entity already sent by said trading entity to at least one of said plurality of auction entities;

20

constructing a plurality of currently successful sets of said offer data of said plurality of auction entities;

determining a value to undercut each said set of currently successful  
25 offers;

selecting a said currently successful set of offers having an optimum value to undercut; and



30990066

-35-

constructing a set of offers of said trading device which minimally undercut said optimum value to undercut currently successful offer set.

9. The method as claimed in claim 5, wherein said step of  
5 constructing a set of offers of said trading device which are lower than said optimum value to undercut currently successful offer set comprises;

for each offer of said optimum offer set, said offer made either by said trading entity or a third party, selecting a lowest third party offer appearing in an  
10 identical auction entity to said offer, which also appears in said optimum offer set; and

generating a trading device offer of an amount equal to said highest third party offer minus a minimum bid decrement.  
15

10. A method of exchanging data between a first electronic trading entity and a plurality of electronic auction entities, said method comprising the steps of:

20 said first trading entity monitoring data displayed by said plurality of auction entities;

said first trading entity extracting from said monitored data, data describing prices of individual items to be traded;  
25

said first trading entity determining from said price data an optimum set of trading data for sending to said plurality of auction entities; and

said first trading entity communicating said trading data to said plurality of  
30 auction entities.

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30990066

-36-

**Abstract****AUCTION METHOD AND APPARATUS FOR ELECTRONIC COMMERCE**

An electronic trading entity (400) comprises a computing device having a processor, memory, user interface, and communications functionality, the computer entity comprising an algorithm (700) for monitoring a plurality of auction entities remotely over the internet; a algorithm (701) for calculating bids and for purchase of goods/services with a plurality of auction entities; an algorithm (702) for selecting individual auction entities with which to trade; a web browser (506) to allow a user to search for auction entities and add these to a stored list of entities at the bidding device. Bid calculation and offer calculation algorithms operate to optimally place bids at a lowest overall cost to the trading entity, or place offers with the object of obtaining the highest overall monetary figures for goods or services for sale by the trading entity.

Fig. 7

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1/20

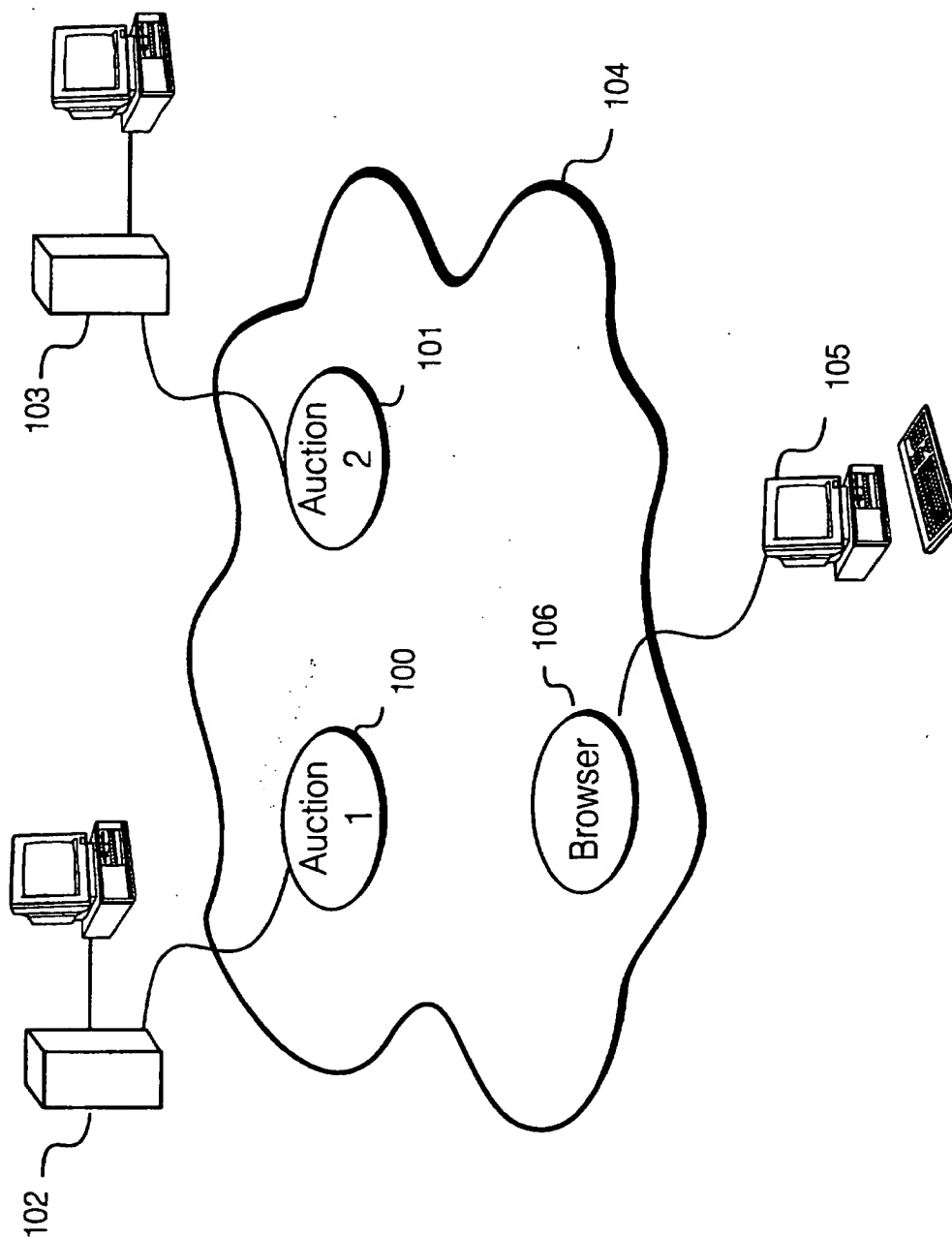


Fig. 1  
(prior art)

2/20

Outstanding 2.88 Carat Round Diamond VVS2 I
Item #125854960

Jewelry, Gemstones:Gemstones:General

Currently	\$11,500.00 (reserve not yet met)	First bid	\$5,000.00
Quantity	1	# of bids	15 (bid history)
Desc.	Time left	Location	(with emails)
	4 days, 20 hours +	Portland, OR	
	Started	(mail this auction to a friend)	
	07/01/99, 23:37:24 PDT	(request a gift alert)	
	Ends	07/11/99, 23:37:24 PDT	
		Featured Category Auction	
Bid!	Seller	Hendrik (206)	
		(view comments in seller's Feedback profile) (view seller's other auctions)	
		(ask seller a question)	
High bid	Manolod (Q)		
Payment	Money Order/Cashiers Checks, personal Checks, See item description for payment methods accepted		
Shipping	Seller ships internationally, See item description for shipping charges		

Seller assumes all responsibility for listing this item. You should contact the seller to resolve any questions before bidding. Currency is dollar (\$) unless otherwise noted.

Fig. 2  
(Prior Art)

3/20

Bidding

Outstanding 2.88 Carat Round Diamond WVS2 I (Item #125854960)  
eBay item 125854960 (EN...:24 PDT) - Outstanding 2.88 Carat Round Diamond WVS2

Current bid

\$11,500.00

Bid increment

\$100.00

Minimum bid

\$11,600.00

Registration is required in order to bid. Find out how to become a registered user. It's fast and it's free!

To finalize your bid, you will need to submit your User ID and Password in the next step. You will not be asked to enter your User ID and Password anywhere on this page.

Current minimum bid is \$11,600.00

Review bid

Your maximum bid.

Please type only numerals and the decimal symbol (.) if required. Do not include currency symbols such as \$.

Binding contract.

Placing a bid is a binding contract in many states. Do not bid unless you intend to buy this item at the amount of your bid.

Proxy bidding for all bids

Please bid the maximum amount you are willing to pay for this item. Your maximum amount will be kept secret; we will bid on your behalf as necessary by increasing your bid by the current bid increment up until your maximum is reached. This saves you the trouble of having to keep track of the auction as it proceeds and prevents you from being outbid at the last minute unless your spending limit is exceeded. (See an example of proxy bidding). Also, in case of a tie for high bidder, earlier bids take precedence. And, keep in mind that you cannot reduce your maximum bid at a later date. Unless otherwise noted, bids are in U.S. dollars.

If you have bid on this item before, note that your new bid must be greater than your previous bid.

Fig. 3 (Prior Art)

4/20

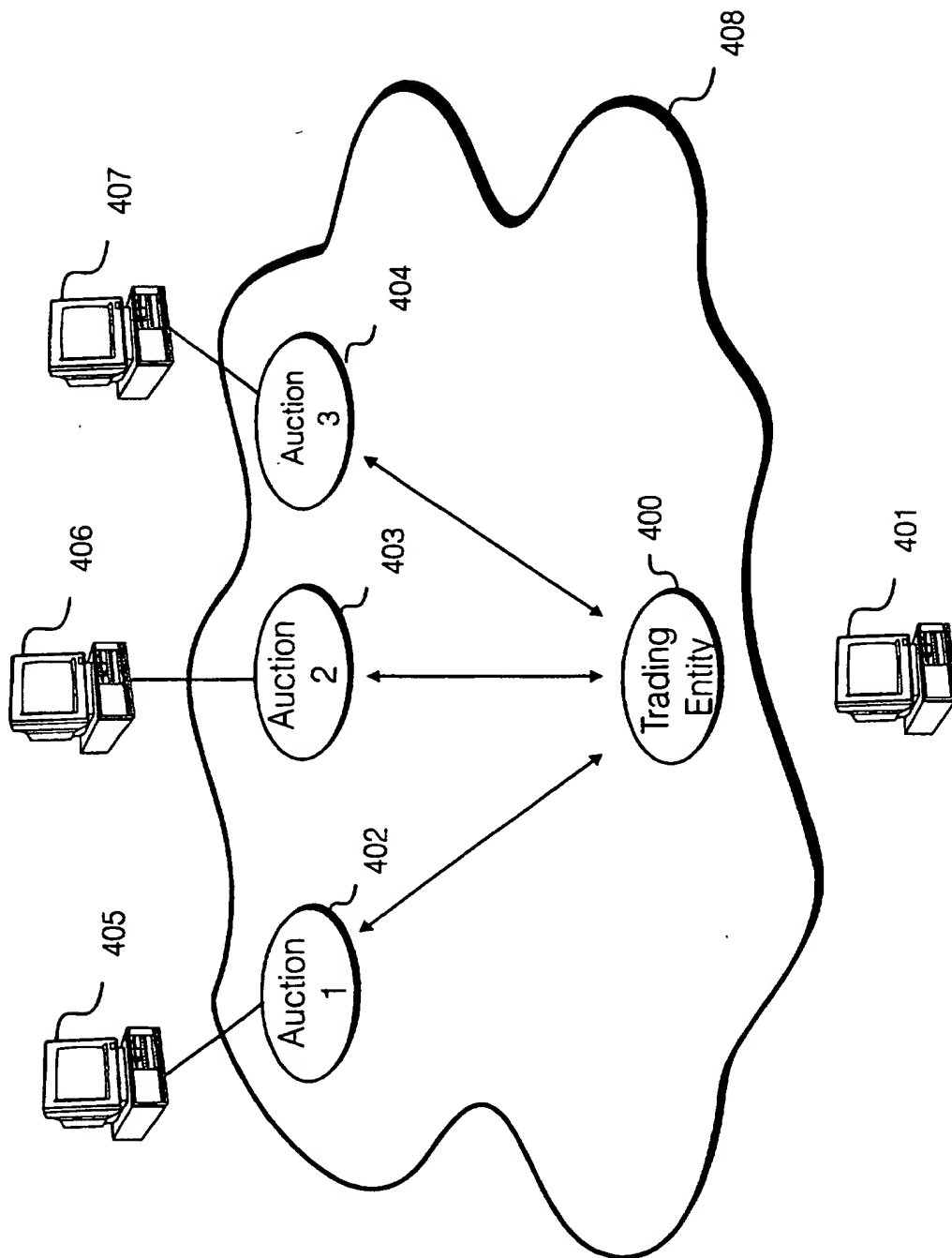


Fig. 4



5/20

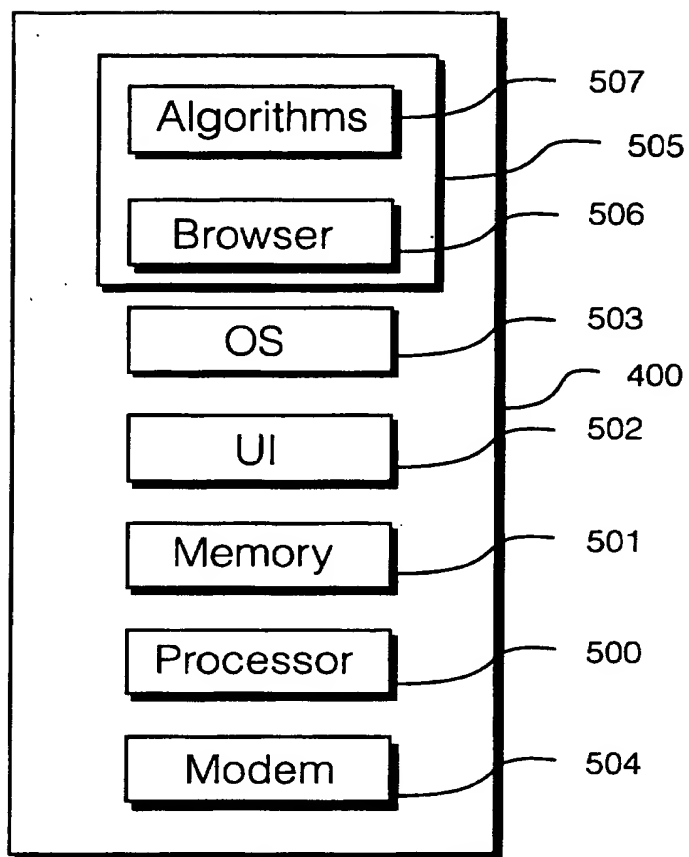


Fig. 5

6/20

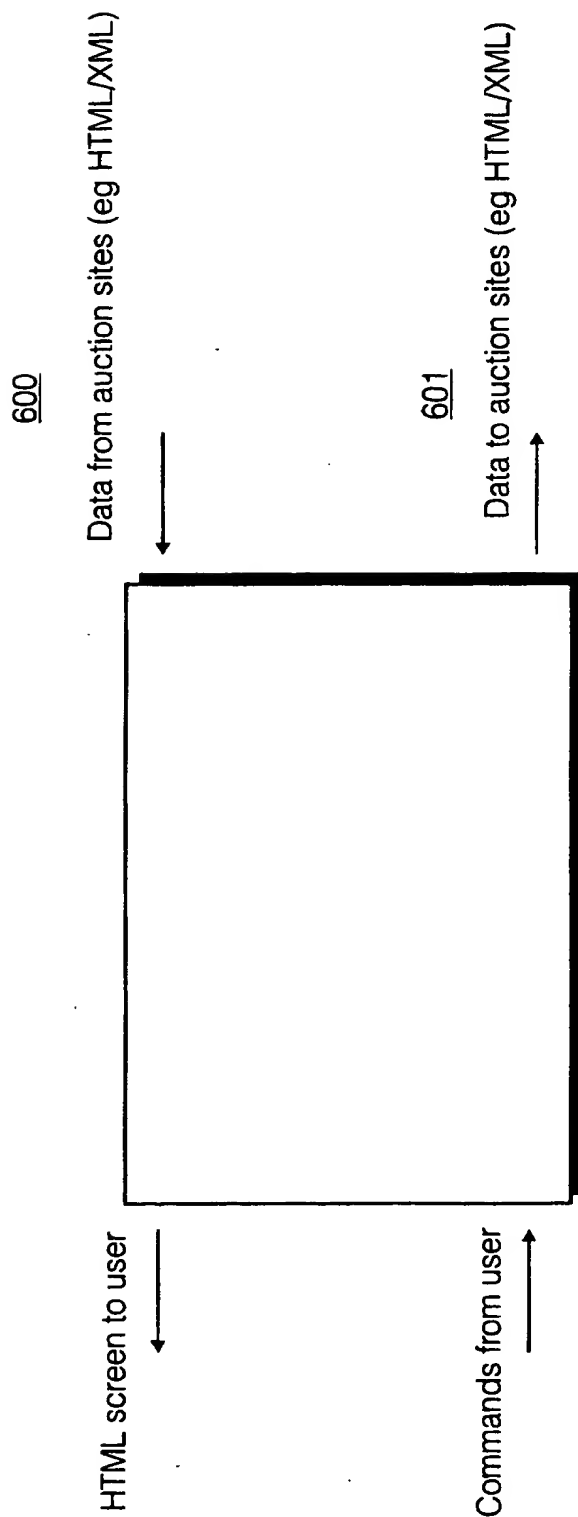


Fig. 6

7/20

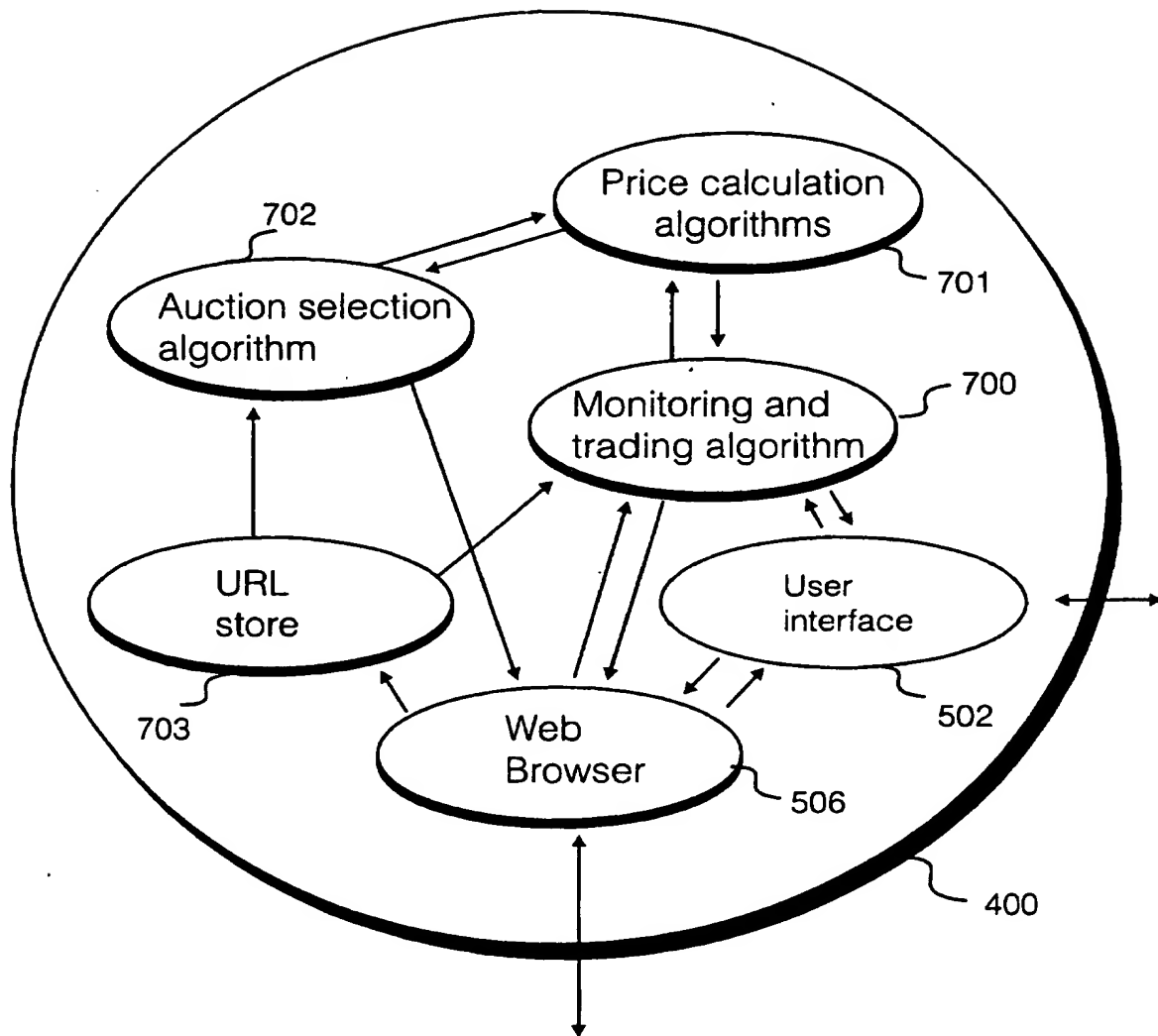


Fig. 7

8/20

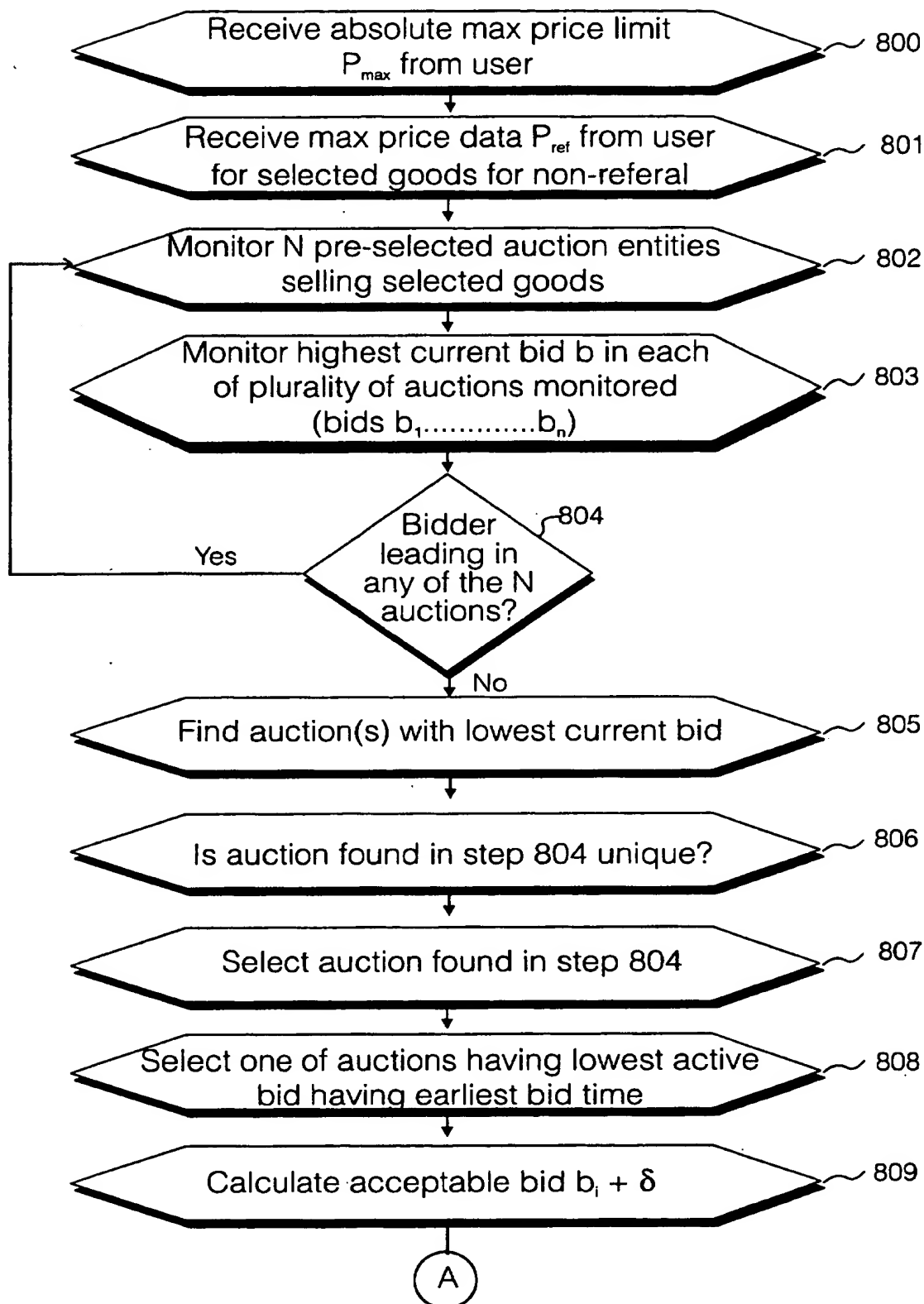


Fig. 8

9/20

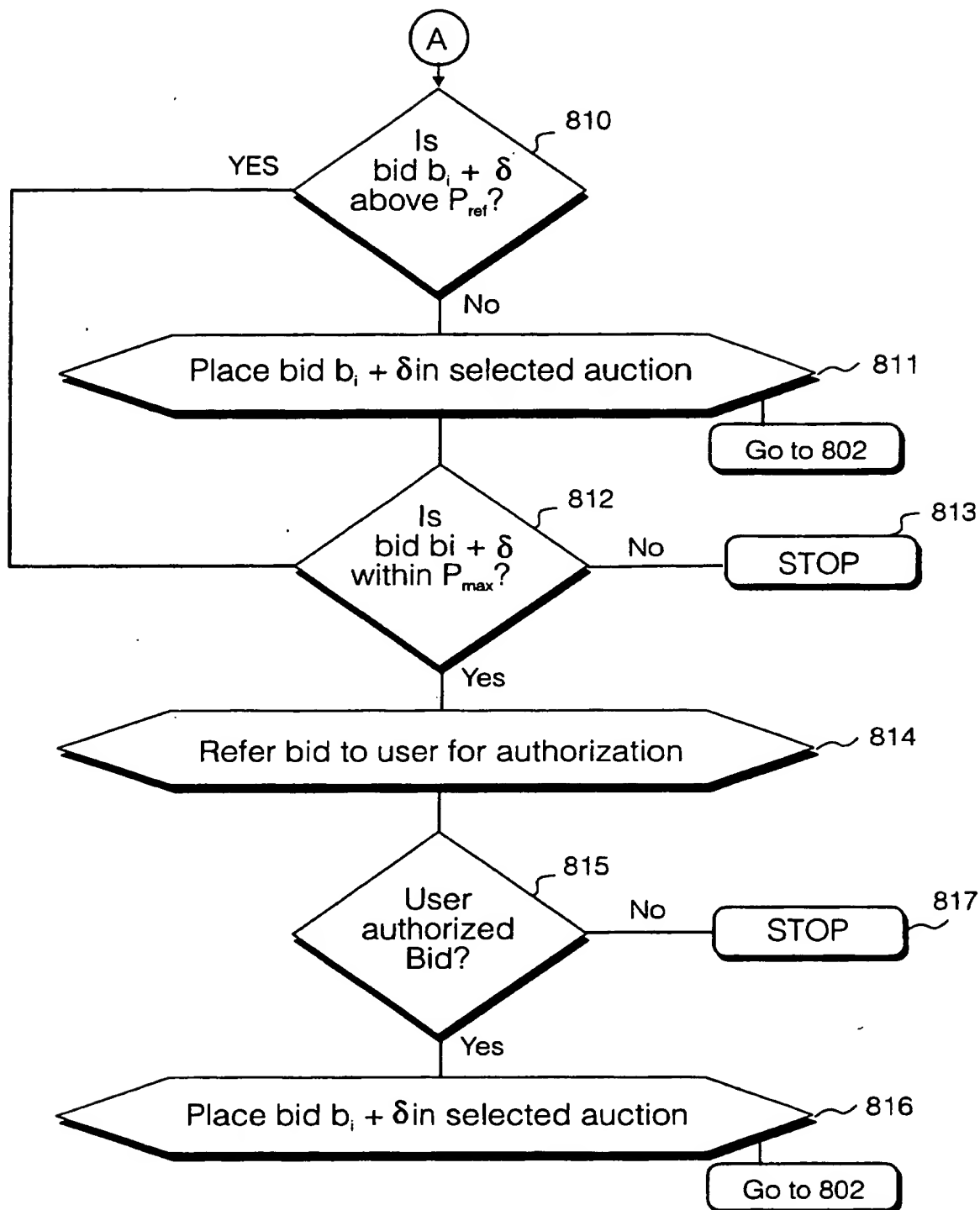


Fig. 8b

10/20

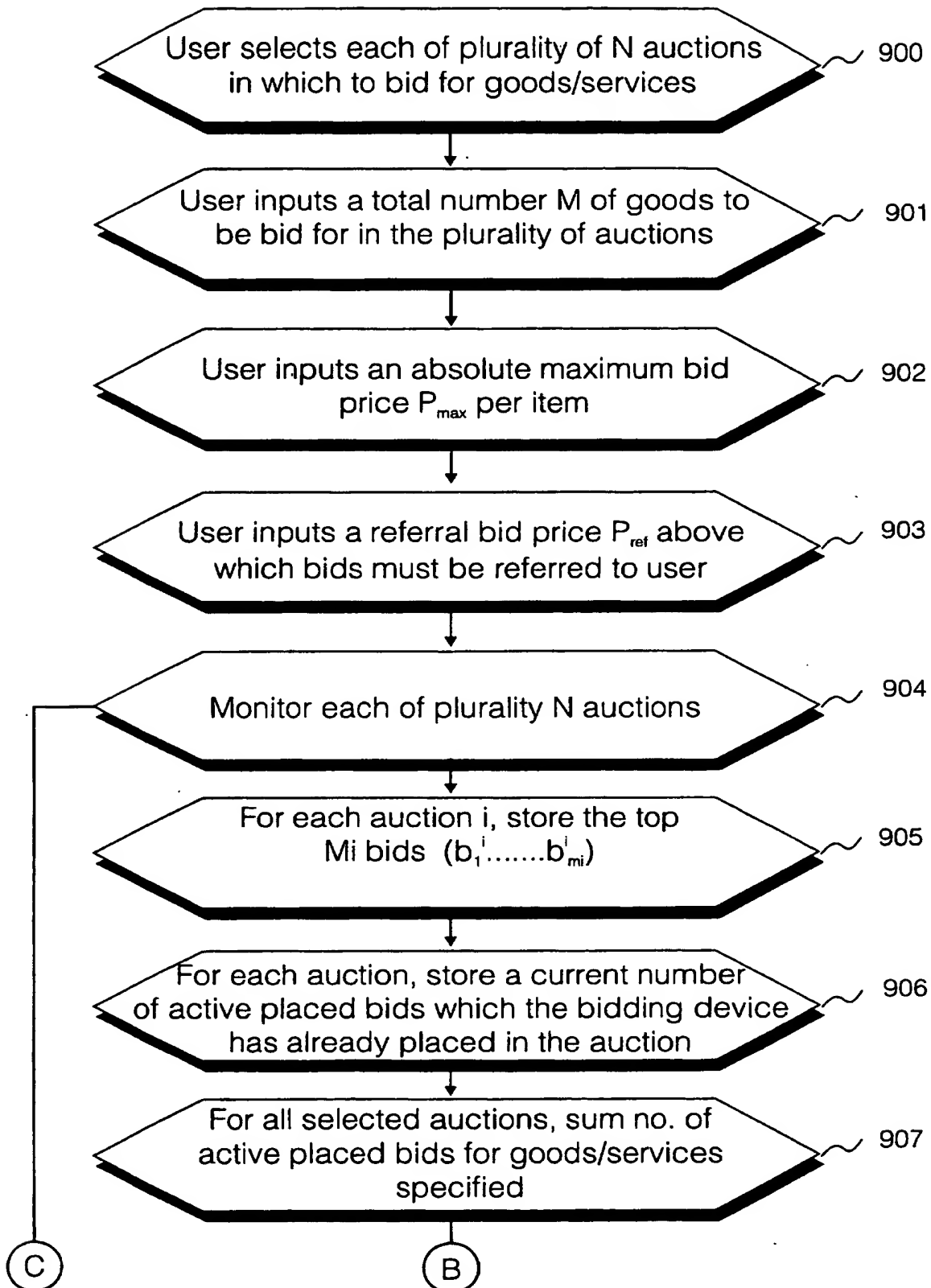


Fig. 9

11/20

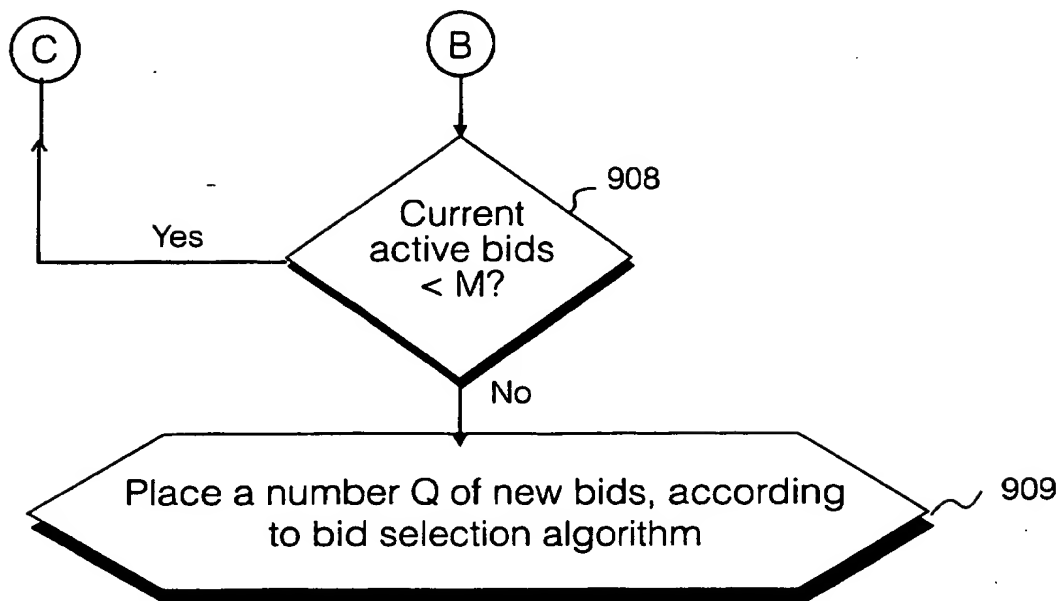


Fig. 9b

12/20

1000	1001	1002	1003	1004	1005	1006	1007	1008
Bid identifier	Auction ID	Lot No./ID	Goods/Services	Lot size (Units)	Amount bid (US Dollars)	Time of bid (GMT)	Close time (GMT)	Status
1278/99	372.278.596	1501	10k register	10,000	25 USD	10/2/99 -12.42	10/7/99 -17.00	Active
1279/99	281.502.7938	1720	10k register	10,000	27 Euro	10/2/99 -12.44	10/2/99 -18.00	Active
1280/99	282.702.6229	10	Dell monitor Model 1010	25	2500 USD	10/2/99 -12.50	10/2/99 -22.00	Active
1281/99	283.701.642	68	HP monitor 2050	25	3000 Euro	10/2/99 -12.59	11/2/99 -12.00	Active
1282/99	7298.73.25	1050231	Dell monitor	10	1100 Euro	10/2/99 -13.02	11/2/99 -12.00	Active
1283/99	25.9827.532	98710	Dell monitor	5	600 USD	10/2/99 -13.04		Active
1284/99	123.789.258	10/A99G	Lexus 200 GLS Automobile	1	25,001 USD	10/2/99 -15.02	15/2/99 -17.00	Outbid
1285/99	372.278.596	1502	10k register	10,000	27 USD	10/2/99 -15.03	10/2/99 -17.00	Active
1286/99	372.278.596	1503	10k register	10,000	28 USD	10/2/99 -15.03	10/2/99 -17.00	Active
1287/99	281.502.7938	1724	10k register	10,000	28 USD	10/2/99 -15.04	10/2/99 -18.00	Active

Fig. 10



13/20

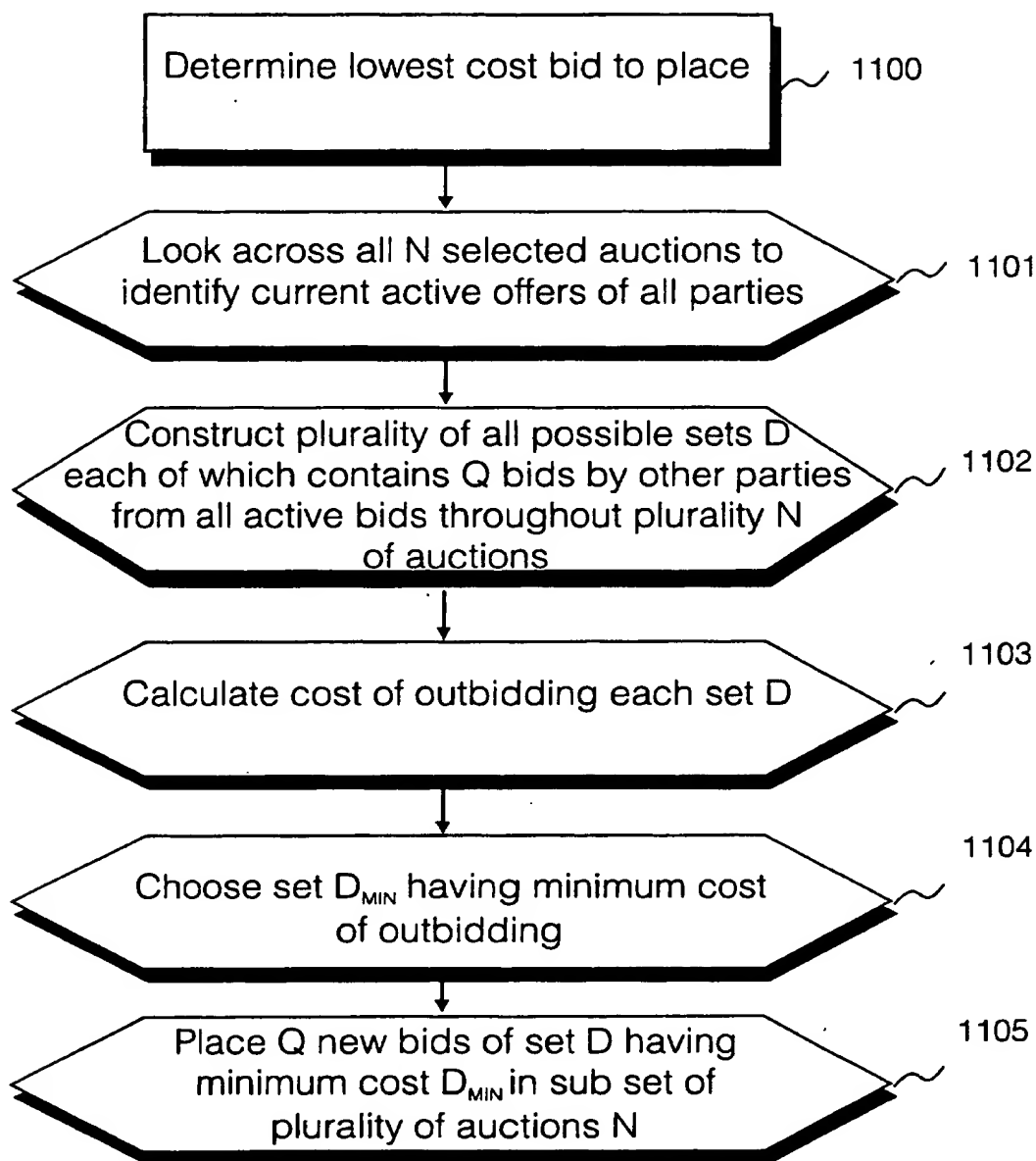


Fig. 11

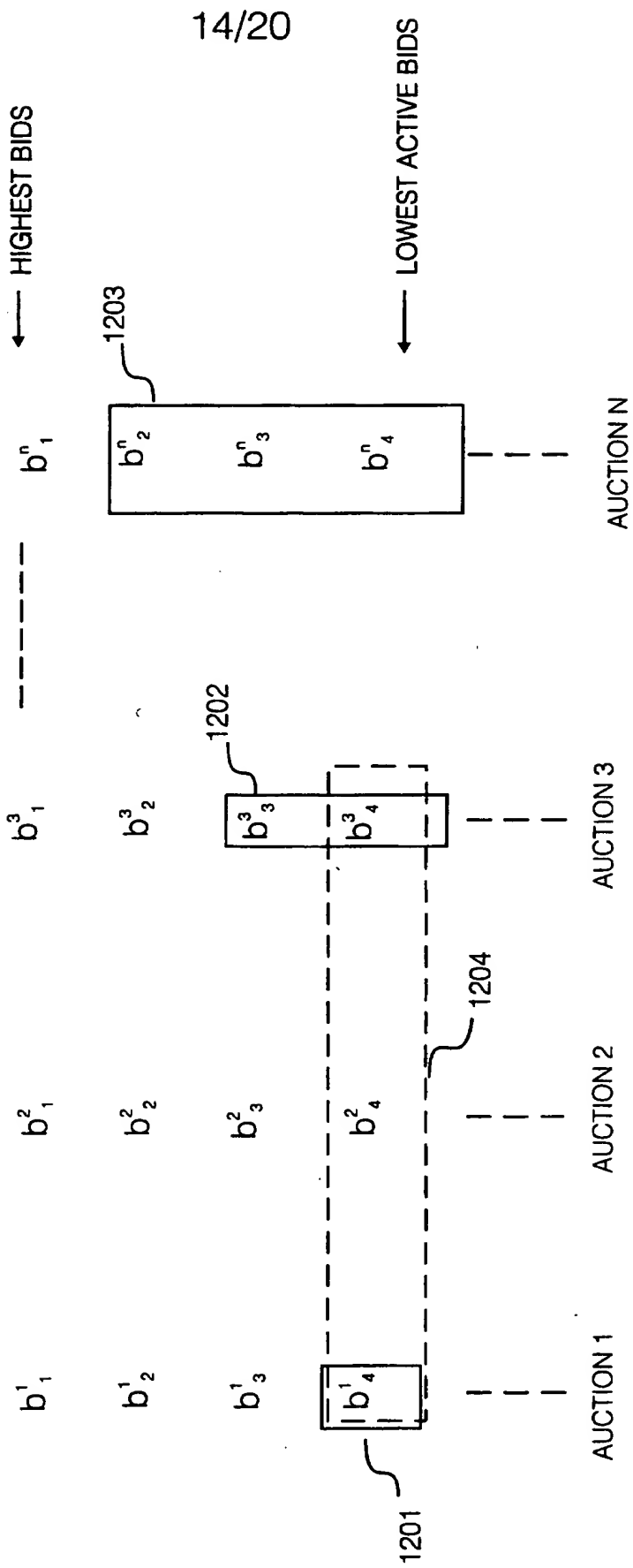


Fig. 12

15/20

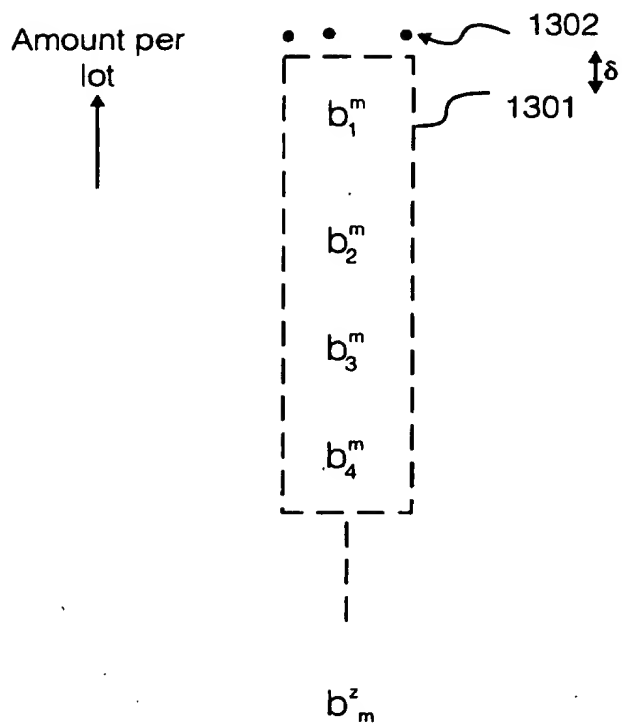


Fig. 13

16/20

## ACTIVE BIDS

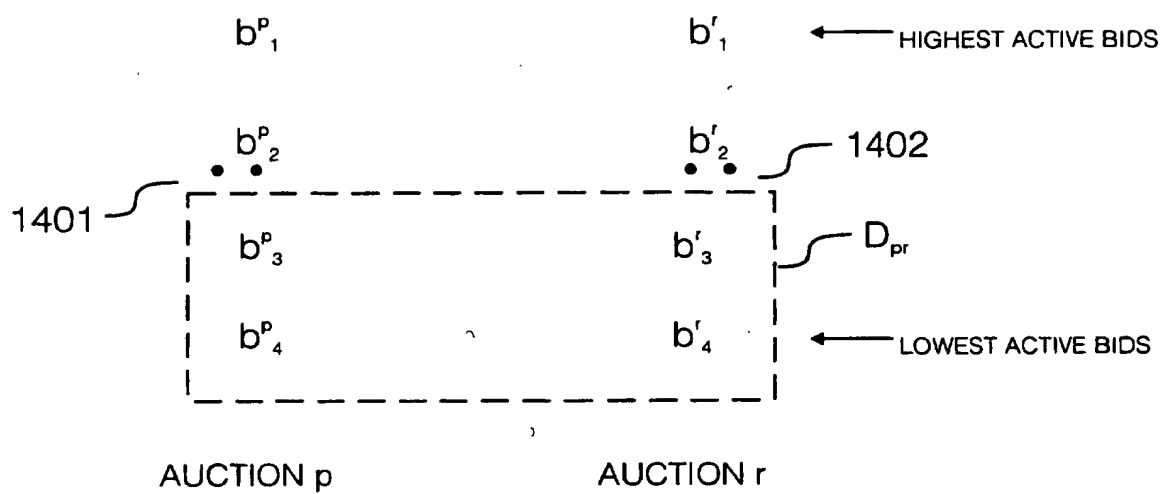


Fig. 14

17/20

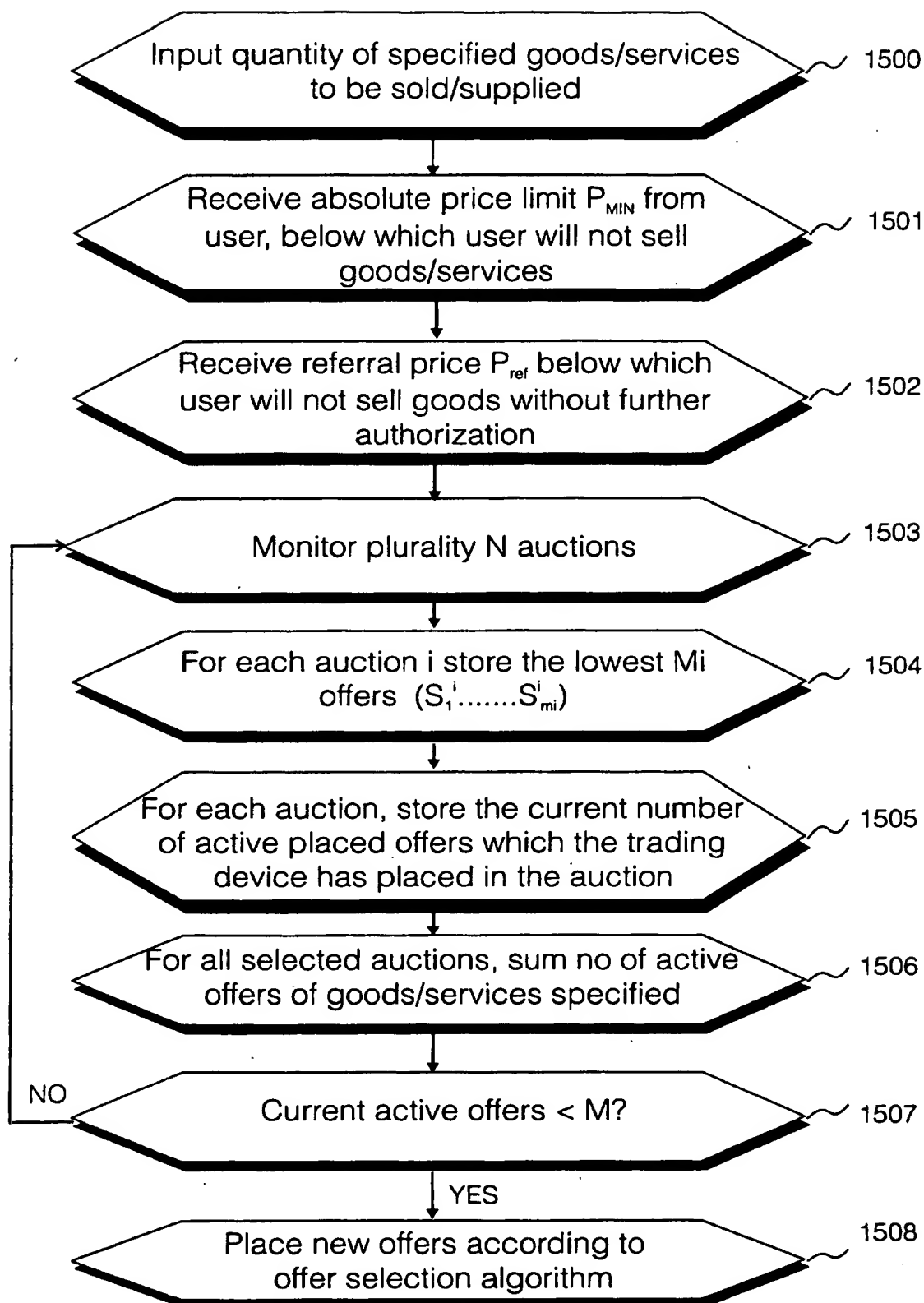


Fig. 15

18/20

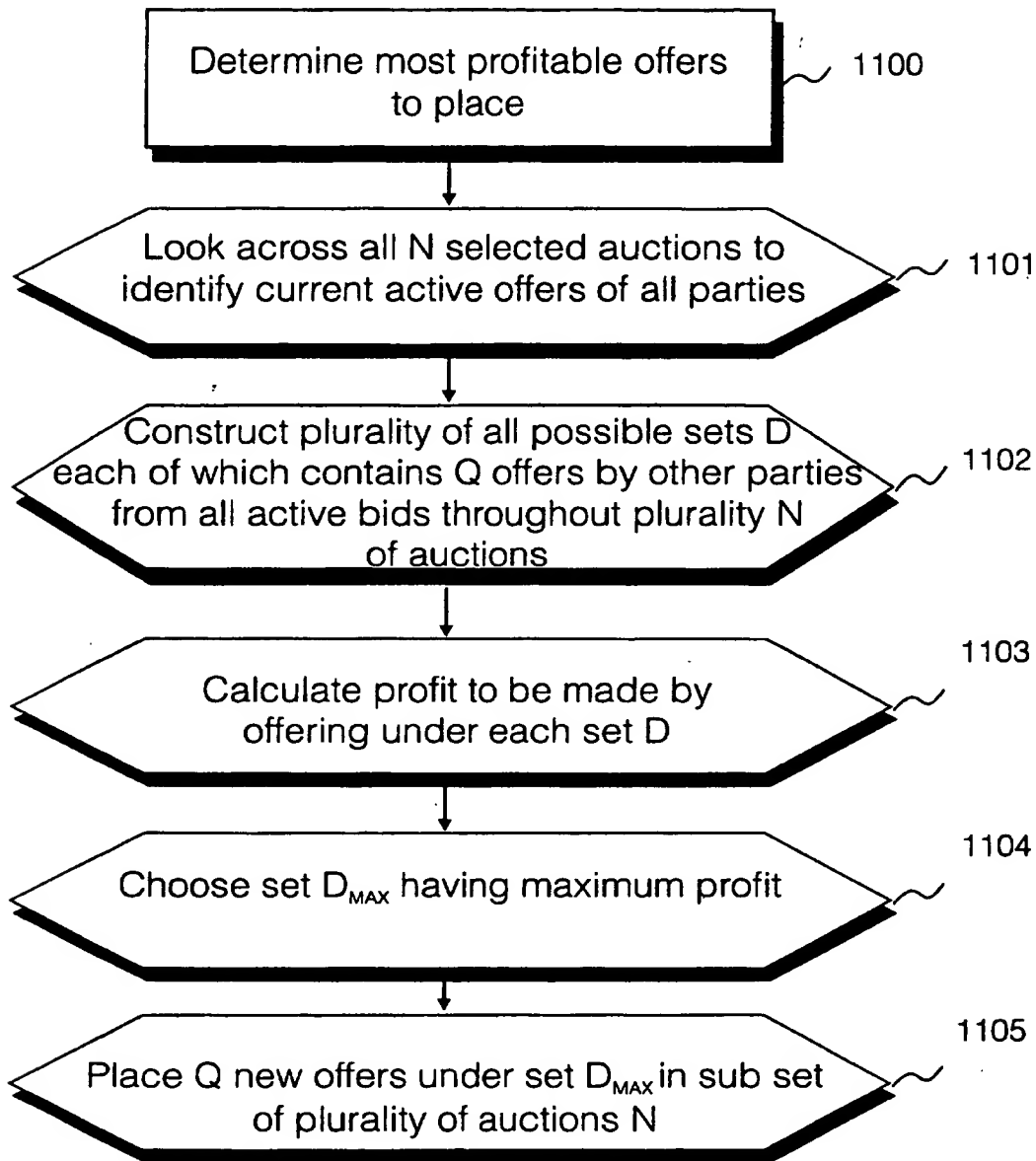


Fig. 16

19/20

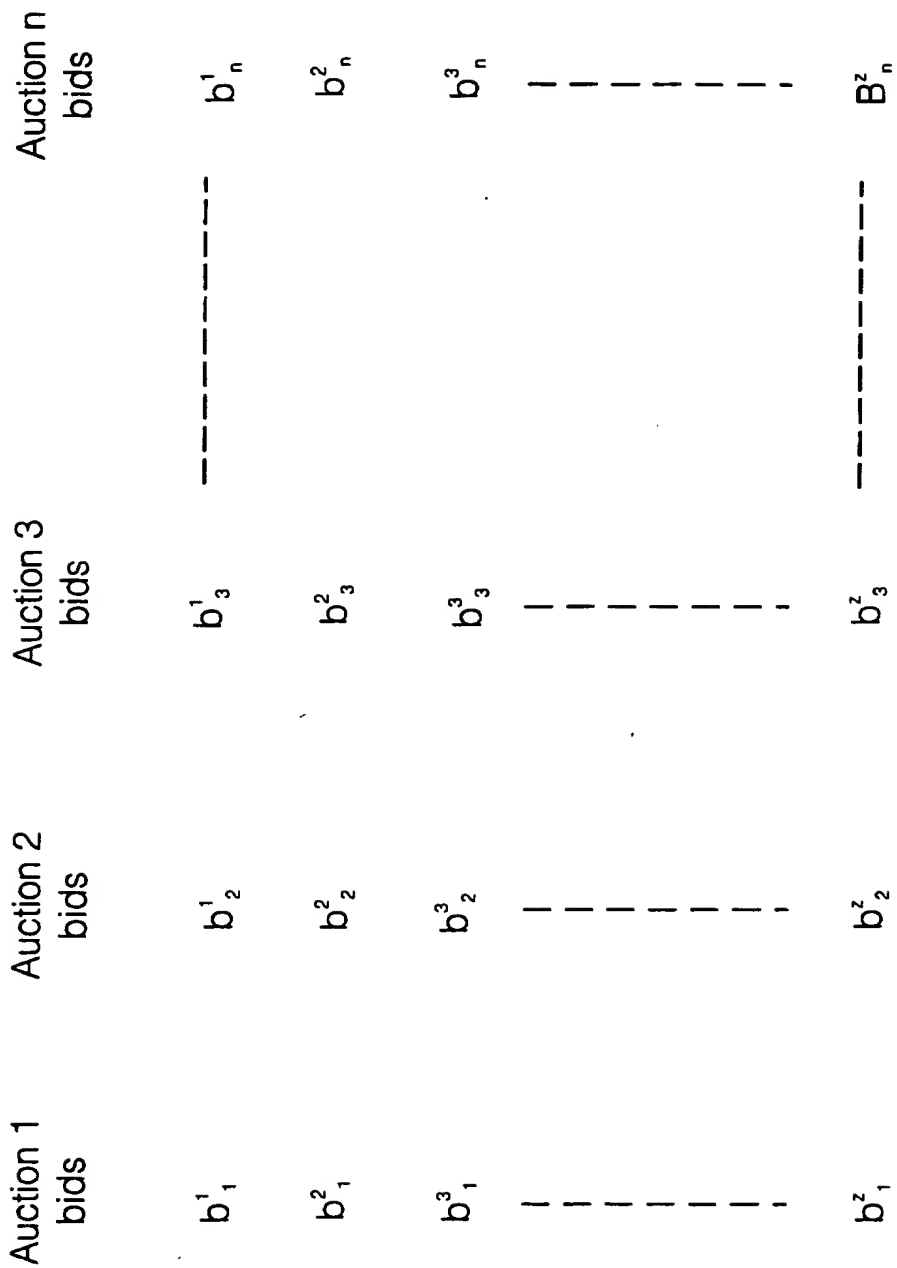


Fig. 17

20/20

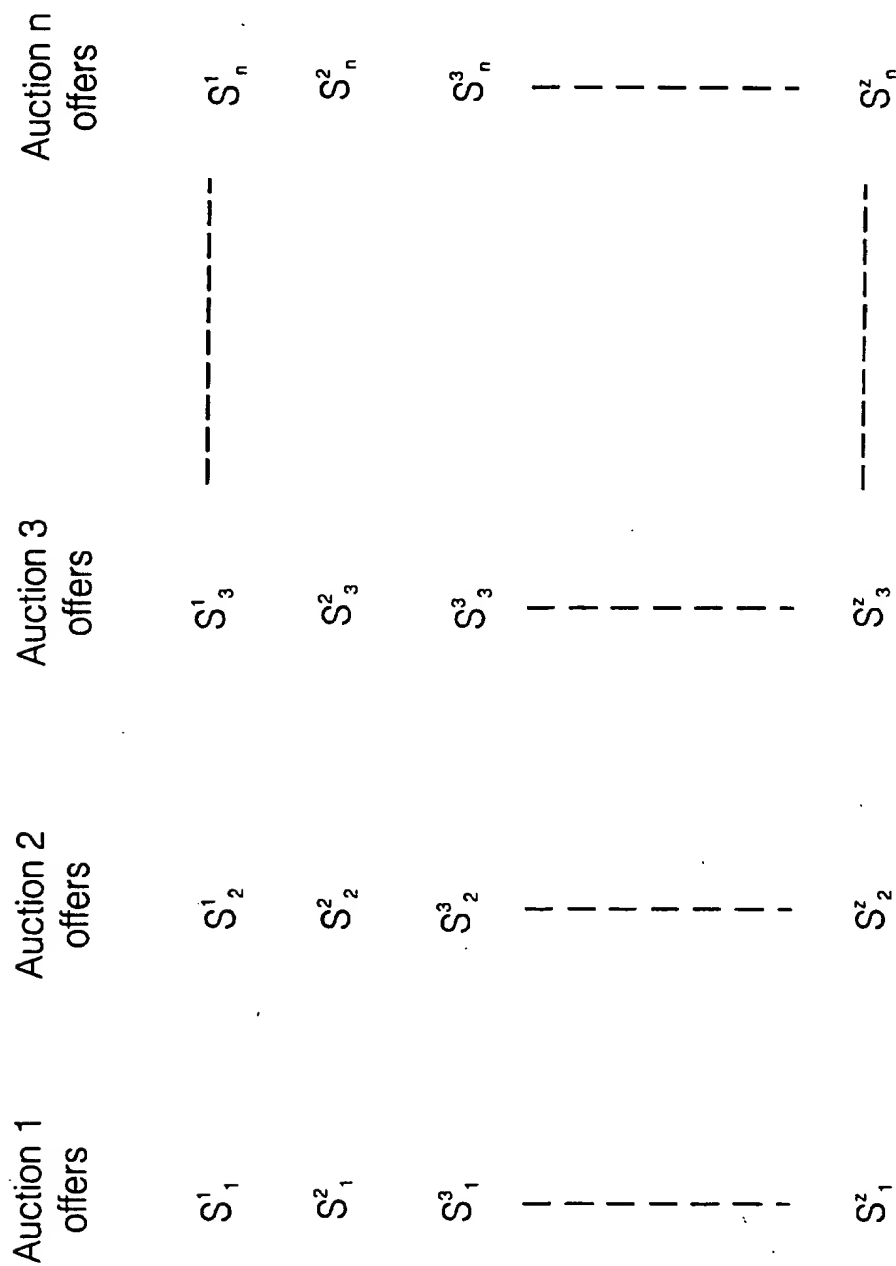


Fig. 18





DOCUMENTS CONSIDERED TO BE RELEVANT									
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int.Cl.7)						
X	WURMAN P R ET AL.: "The Michigan Internet AuctionBot: a configurable auction server for human and software agents" PROCEEDINGS OF THE 2ND INT. CONF. ON AUTONOMOUS AGENTS, 9 - 13 May 1998, pages 301-308, XP002131462 * abstract * * figure 2 * * page 303, right-hand column, line 23 - page 304, left-hand column, line 22 * * page 304, left-hand column, line 33 - page 304, right-hand column, line 11 *	1-10	G06F17/60						
X	MAES P ET AL: "Agents That Buy and Sell" COMMUNICATIONS OF THE ACM, vol. 42, no. 3, March 1999 (1999-03), pages 81-91, XP002131463 grey underlayed section "Automated Negotiation" on pp 84-85	1,5,10							
A	TSVETOVATYY M ET AL: "MAGMA: an agent-based virtual market for electronic commerce" APPLIED ARTIFICIAL INTELLIGENCE, 22 - 24 April 1996, XP002079276 * the whole document *	1,5,10	TECHNICAL FIELDS SEARCHED (Int.Cl.7) G06F						
A	US 5 715 402 A (POPOLO CARL A) 3 February 1998 (1998-02-03) * abstract * * figures 3,6-8 * * column 1, line 46 - column 2, line 12 * -/-	1,5,10							
The present search report has been drawn up for all claims									
Place of search BERLIN		Date of completion of the search 24 February 2000	Examiner Schmidt, A						
<table border="0"><tr><td>CATEGORY OF CITED DOCUMENTS</td><td>T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons ----- &amp; : member of the same patent family, corresponding document</td></tr><tr><td>X : particularly relevant if taken alone Y : particularly relevant if combined with another document of the same category A : technological background O : non-written disclosure P : intermediate document</td><td></td><td></td><td></td></tr></table>				CATEGORY OF CITED DOCUMENTS	T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons ----- & : member of the same patent family, corresponding document	X : particularly relevant if taken alone Y : particularly relevant if combined with another document of the same category A : technological background O : non-written disclosure P : intermediate document			
CATEGORY OF CITED DOCUMENTS	T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons ----- & : member of the same patent family, corresponding document								
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DOCUMENTS CONSIDERED TO BE RELEVANT			
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (IntCl.7)
A	WO 98 34187 A (GROEFTEHAUGE MARTIN ;AUTOCOM APS (DK); GROEFTEHAUGE PETER (DK)) 6 August 1998 (1998-08-06) * abstract * * page 1, line 1 - page 5, line 5 * * figure 1 *  ————	1,5,10	
			TECHNICAL FIELDS SEARCHED (IntCl.7)
The present search report has been drawn up for all claims			
Place of search <b>BERLIN</b>		Date of completion of the search <b>24 February 2000</b>	Examiner <b>Schmidt, A</b>
<b>CATEGORY OF CITED DOCUMENTS</b> X : particularly relevant if taken alone Y : particularly relevant if combined with another document of the same category A : technological background O : non-written disclosure P : intermediate document  T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons  & : member of the same patent family, corresponding document			

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**EP 99 30 7307**

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The European Patent Office is in no way liable for these particulars which are merely given for the purpose of information.

**24-02-2000**

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